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# Computer Interfaces

David Vernon

# Computer Interfaces

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- ◆ The goal of this course is to introduce hardware and software design techniques and issues for interfacing computers and peripheral devices

# Computer Interfaces

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## ◆ Course Contents

- ❑ Digital data communication standards – serial communications
  - Devices
  - RS232
  - RS422
  - Handshaking
  - Implementation of RS232 on a PC
- ❑ Universal Serial Bus (USB)
  - USB standards
  - Types and elements of USB transfers
  - Development procedure for USB applications

# Computer Interfaces

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## ◆ Course Contents

### □ Parallel Communications

- General Purpose Interface Bus (GPIB)
- GPIB signals and lines
- Handshaking and interface management
- Implementation of a GPIB on a PC

# Computer Interfaces

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## ◆ Course Contents

### □ Digital and Analogue Interfacing

- Digital Interfacing
  - o Digital I/O ports
  - o Interfacing external signals to digital I/O ports
  - o Optical isolation
- Analogue Interfacing
  - o Revision of A/D and D/A conversion techniques
  - o Multiplexing
  - o Analogue I/O cards
  - o Data acquisition and control using a PC

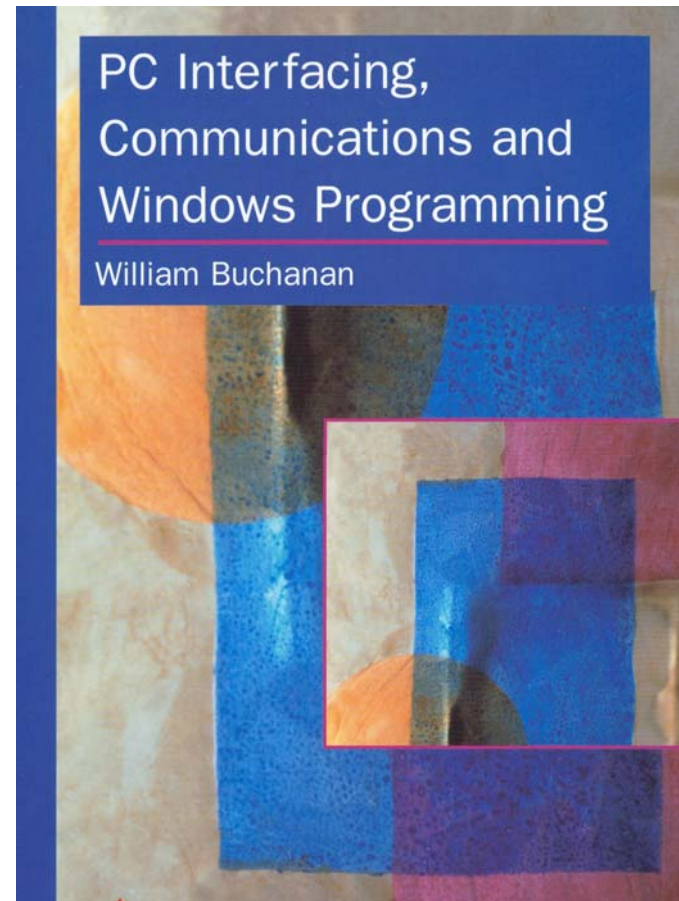
# Computer Interfaces

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## ◆ Course Texts:

*PC Interfacing  
Communications and  
Windows Programming,*  
William Buchanan, Addison  
Wesley

*Microcomputer Interfacing  
and Applications,*  
M. A. Mustafa, Newnes



# Computer Interfaces

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## ◆ Course Loading: 0.5 Module

*1 Lecture each week  
(weeks 1, 2, 3, ...)*

*1 Tutorial every second week  
(weeks 3, 5, 7, 10, 12, 14)*

*1 Lab demonstration every fourth week  
(weeks 4, 11, 15)*

*1 Assignment, due on the last day of week 15 (22<sup>nd</sup> Dec.)*

# Computer Interfaces

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## Motivation



# Computer Interfaces

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# Computer Interfaces

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## ◆ The need for computer interfacing

- ❑ Advanced control applications need flexible processing power, i.e. computers
- ❑ Control data has to be input and output
  - Input from sensors (speed, acceleration, temperature, ..)
  - Output to actuators (motors, switches, valves, ...)
- ❑ Examples:
  - Robotics
  - Industrial process control
- ❑ Advantages of using computers for Data Acquisition & Control
  - High speed
  - Programming flexibility (compared with hard-wired logic)
  - Mass storage of data
  - Data analysis and visualization
  - Low cost (relatively)

# Computer Interfaces

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## ◆ Issues:

- ❑ Input vs output (direction)
  - Communication between devices
  - Data acquisition
  - Device control
- ❑ Digital vs Analogue data (modality)
  - Signal levels (may be a need for signal conditioning)
  - Analogue to digital conversion (ADC)
  - Digital to analogue conversion (DAC)
- ❑ Serial vs parallel (mechanism)
  - Speed, distance, number of required lines, standard I/F
- ❑ Polled vs interrupt driven (mechanism)
  - Simplicity, processor efficiency

# Computer Interfaces

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## Microprocessor Architecture

# Microprocessor Architecture

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## ◆ History

### ❑ Early '70s: Intel 4004

- 4 bit of data (nibble)
- 2000 transistors
- 46 instructions
- 4kB program code
- 1kB data

### ❑ 1974 - : 8008, 8080, 8085

- 8008: 14-bit address space (16kB memory)
- 8080: 16-bit address space (64kB memory)

# Microprocessor Architecture

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## ◆ History

- ❑ Early 80s: 8086 – Major revolution in processing power
  - 16-bit data bus
  - 20 bit address bus (1MB memory)
  - PC & PC-XT
  - 8088: Multiplexed 8-bit data bus
  - 80286: enhance version (PC AT)

# Microprocessor Architecture

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## ◆ History

- ❑ 1985: Intel 80386
  - 32 bit data bus
  - 32-bit address bus (4GB memory)
- ❑ 1986: Intel 80486
  - Memory cache
  - On-chip maths co-processor

# Microprocessor Architecture

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## ◆ History

### ❑ 1990s: Pentium (P-5)

- 64-bit superscalar architecture
- MISD instruction pipeline (can execute more than one instruction at a time)
- 64-bit databus
- 32-bit address bus

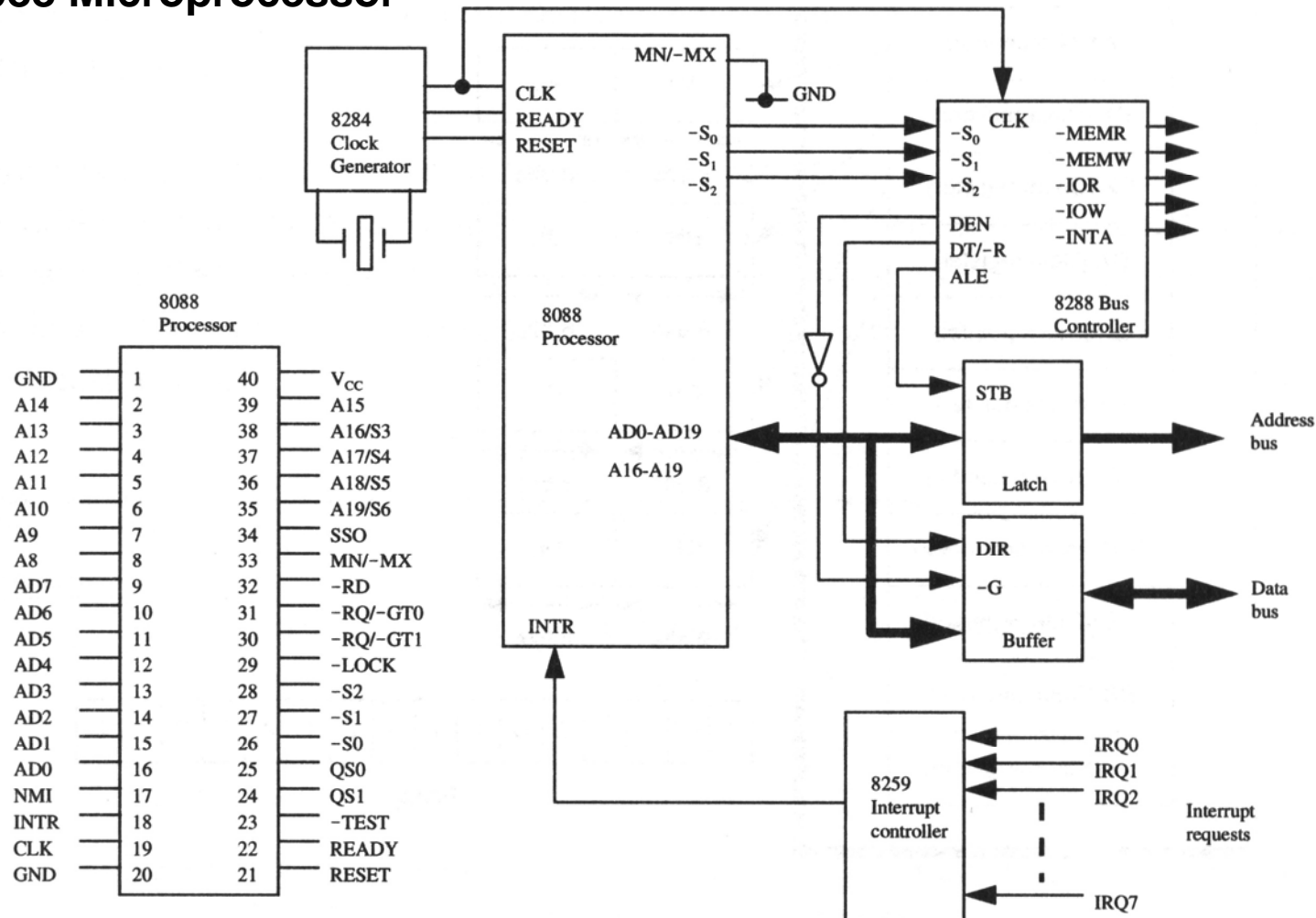
### ❑ 2000s: Pentium II (P-6)

- Up to 4 processors on the same bus
- Single bit error detection and correction on the data-bus
- Multiple bit error detection



# Microprocessor Architecture

## 8088 Microprocessor



# Microprocessor Architecture

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## ◆ 8088

- ❑ 8288 bus controller generates signals based on the 8088 status lines  $\overline{S0}$ ,  $\overline{S1}$ ,  $\overline{S2}$ :
- ❑ Main control signals:
  - $\overline{IOR}$  (I/O Read): processor is reading from the content of the address which is on the **I/O** bus
  - $\overline{IOW}$  (I/O Write): processor is writing the contents of the data bus to the address on the **I/O** bus
  - $\overline{MEMR}$  (memory read): processor is reading from the contents of the address which is on the **address** bus
  - $\overline{MEMW}$  (memory write): processor is writing the contents of the data bus to the address which is on the address bus
  - $\overline{INTA}$  (interrupt acknowledge): used by the processor to acknowledge an interrupt (sent via the 8259 interrupt controller)

# Microprocessor Architecture

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## ◆ 8088

### □ Two address spaces:

- Memory space
- I/O space

### □ Registers

- 4 general purpose registers

AX Accumulator (all I/O operations & some arithmetic)

BX base register (can be used as an address register)

CX count register (e.g. loops)

DX data register (some I/O and when multiplying & dividing)

- 4 addressing registers

SI source index (extended addressing commands)

DI destination index (used in some addressing modes)

BP base pointer

SP stack pointer

# Microprocessor Architecture

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## ◆ 8088

### □ Registers

- Status registers registers

  - IP instruction pointer (contains address of next instruction to be executed)

  - F flag register (with bits set depending on conditions or results of operations, e.g. is the result negative)

- Segment registers (16 bits – 64kB)

  - CS register: code segment; defines the memory location where the code/instructions are stored

  - DS register: data segment; defined where the data from the program will be stored

  - SS register: stack segment; location of the stack

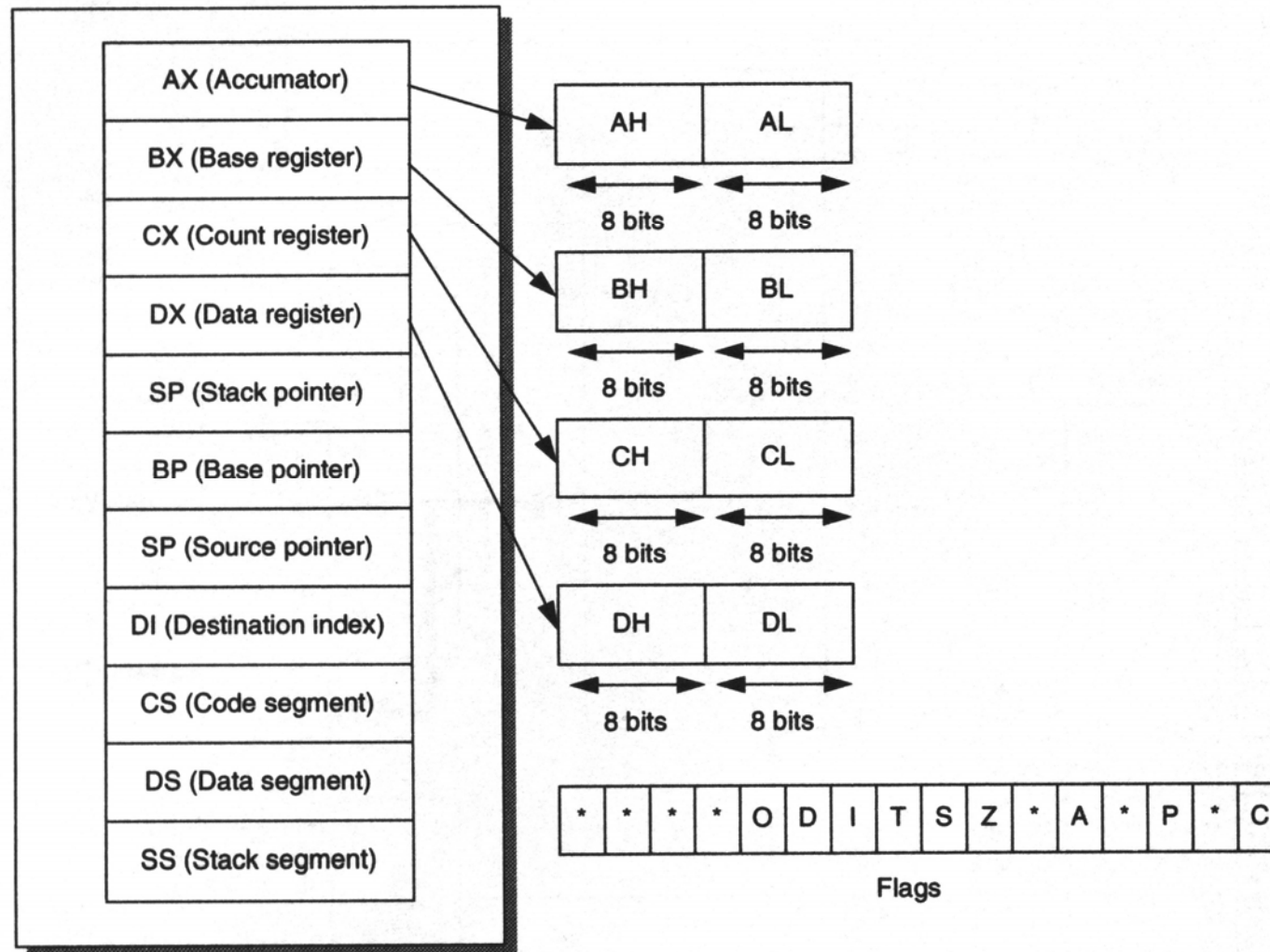
  - ES register: extra segment

- All addresses are defined with respect to the segment registers

- 8086 has a *segmented memory architecture*

# Microprocessor Architecture

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# Microprocessor Architecture

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## ◆ Memory segmentation

- ☐ 8086 has a 20-bit address space (1MB)
- ☐ But can only directly address 64kB (16-bit address space)
- ☐ I.e. it addresses the 1MB in chunks or segments
- ☐ A segmented memory address location is identified with a segment and an offset address (this is the logical address)

`segment : offset`

16 bits : 16 bits

- ☐ The actual physical address is calculated by shifting the segment address 4 bits to the left and adding the offset!
- ☐ What's the actual physical address in the following example?

# Microprocessor Architecture

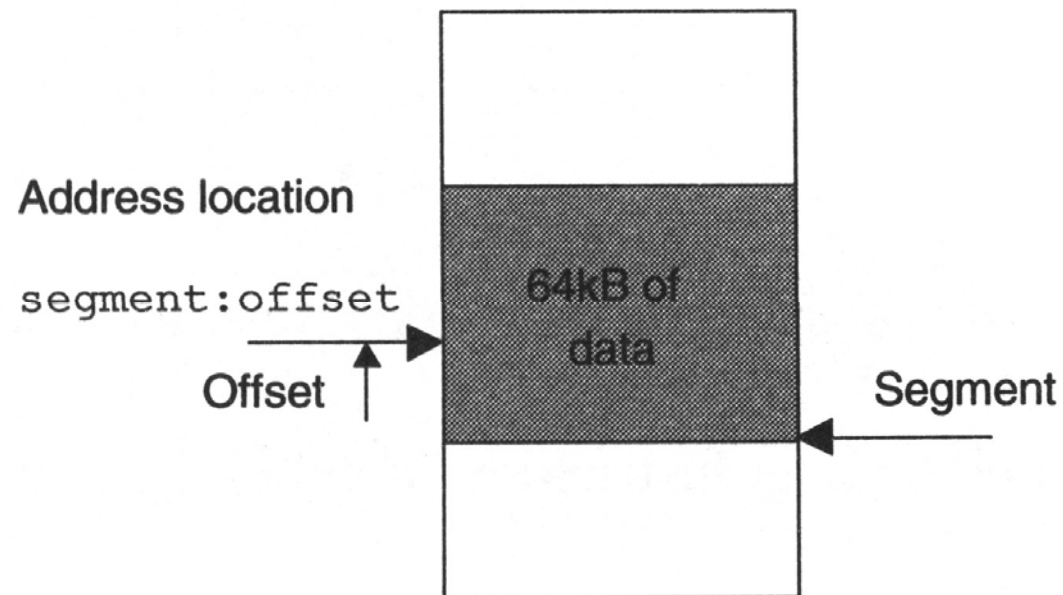
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Segment (2F84):

Offset (0532):

Actual address:

0010	1111	1000	0100	0000
	0000	0101	0011	0010
0010	1111	1101	0111	0010



# Microprocessor Architecture

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## ◆ Memory segmentation

- ❑ In C the address 1234:9876h is specified as 0x12349876

- ❑ Near and far pointers

- A near pointer is a 16-bit pointer (64kB of data)
- A far pointer is a 20-bit pointer (1MB)

```
char far *ptr; /* declare a far pointer */  
ptr = (char far *) 0x1234567; /*initialize a far pointer */
```



# Microprocessor Architecture

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## ◆ Memory mapped I/O vs. Isolated I/O

### ☐ Memory mapped I/O

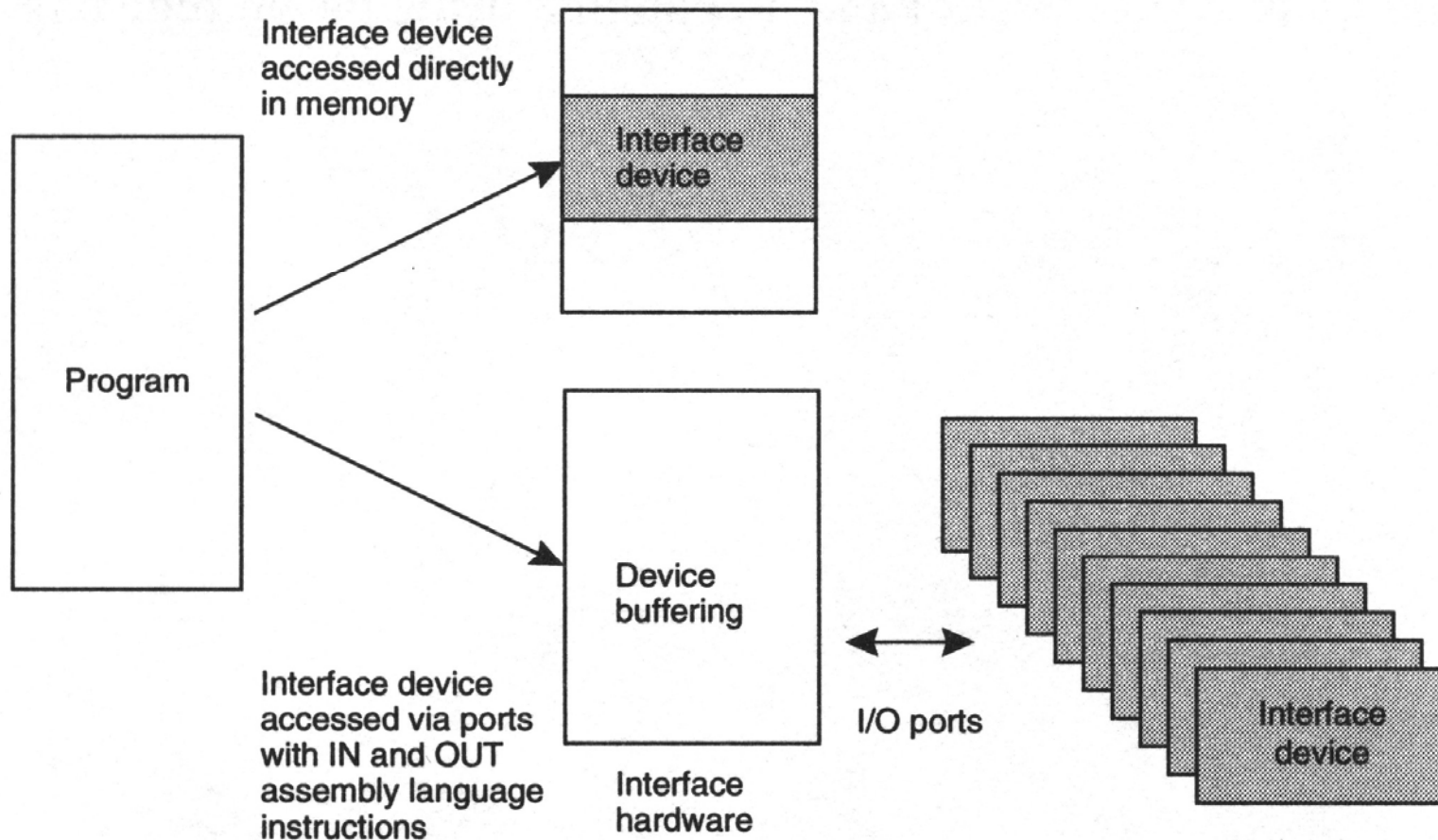
- Devices are mapped into the physical memory space
- Given real addresses on the address bus

### ☐ Isolated I/O

- Devices are mapped into a special isolated memory space
- Accessed via ports which act as a buffer between the processor and the device
- Accessed using the IN and OUT instructions (there are C equivalents)
- Isolated I/O uses 16-bit addressing from 0000h to FFFFh

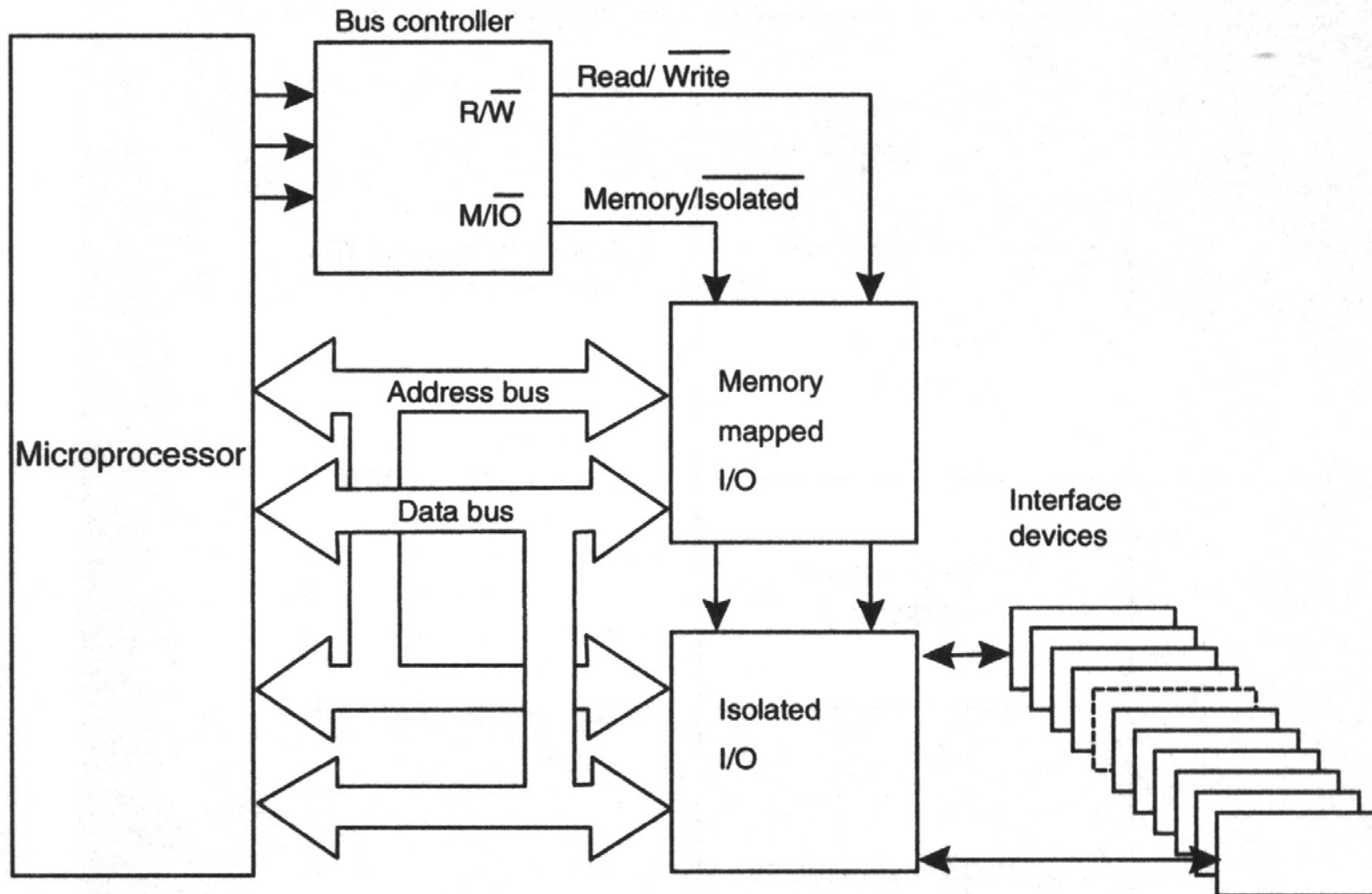
# Microprocessor Architecture

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# Microprocessor Architecture

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# Microprocessor Architecture

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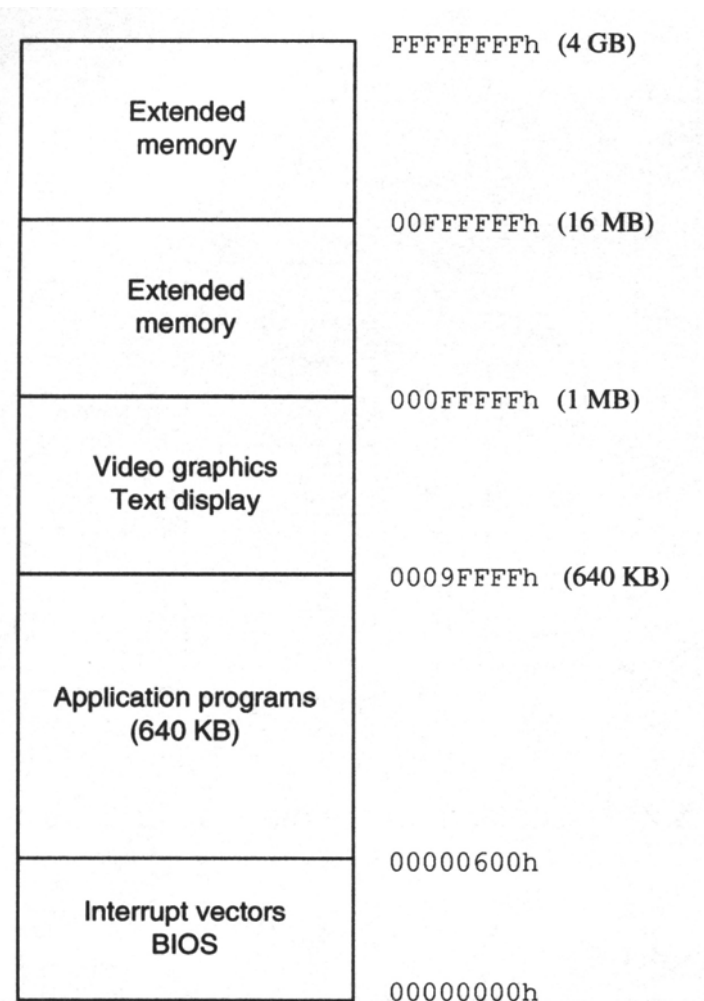
## ◆ 80x86 bus controller signals

### □ $R/\overline{W}$

- Low when data is being written
- High when data is being read
- $M/\overline{IO}$ 
  - o Low when selecting isolated memory
  - o High when selecting normal memory space

# Microprocessor Architecture

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Typical PC Memory Map

# Microprocessor Architecture

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## ◆ Typical isolated I/O Memory Map

000h-01Fh	DMA controller
020h-021h	Programmable interrupt controller (PIC)
040h-05Fh	Counter/Timer
060h-07Fh	Digital I/O
080h-09Fh	DMA controller
0A0h-0BFh	NMI Reset
0C0h-0DFh	DMA controller
0E0h-0FFh	Math co-processor
170h-178h	Hard disk (secondary IDE drive or CD-ROM drive)
1F0h-1F8h	Hard disk (primary IDE drive)
200h-20Fh	Game I/O adapter
210h-217h	Expansion unit
278h-27Fh	Second parallel port (LPT2:)
2F8h-2FFh	Second serial port (COM2:)
300h-31Fh	Prototype card
378h-37Fh	Primary parallel port (LPT1:)
380h-3AFh	SDLC interface
3A0h-3AFh	Primary binary synchronous port
3B0h-3BFh	Graphics adapter
3C0h-3DFh	Graphics adapter
3F0h-3F7h	Floppy disk controller
3F8h-3FFh	Primary serial port (COM1:)

# Microprocessor Architecture

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## ◆ Isolated I/O

### □ Inputting a byte from an I/O port

#### ▪ Borland C

```
unsigned char value;  
value = inportb(PORTADDRESS);
```

Prototyped in `dos.h`

#### ▪ Microsoft C++

```
unsigned char value;  
value = _inp(PORTADDRESS);
```

Prototyped in `conio.h`

# Microprocessor Architecture

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## ◆ Isolated I/O

### □ Inputting a word from an I/O port

#### ▪ Borland C

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unsigned int value;  
value = inport(PORTADDRESS);
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Prototyped in `dos.h`

#### ▪ Microsoft C++

```
unsigned int value;  
value = _inpw(PORTADDRESS);
```

Prototyped in `conio.h`



# Microprocessor Architecture

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## ◆ Isolated I/O

### ❑ Outputting a byte from an I/O port

#### ▪ Borland C

```
unsigned char value;  
outportb(PORTADDRESS, value);
```

Prototyped in `dos.h`

#### ▪ Microsoft C++

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unsigned char value;  
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Prototyped in `conio.h`

# Microprocessor Architecture

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## ◆ Isolated I/O

### □ Outputting a word from an I/O port

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unsigned int value;  
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Prototyped in `dos.h`

#### ▪ Microsoft C++

```
unsigned int value;  
_outw(PORTADDRESS, value);
```

Prototyped in `conio.h`

# Serial Interfaces: RS-232

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- ◆ Serial data transmission is used for digital communication between
  - ❑ Sensors and computers
  - ❑ Computers and computers
  - ❑ Computers and peripheral devices (printer, stylus, mouse, ..)
  - ❑ One of the most widely used communication techniques to interface external equipment
- ◆ Serial communication protocol: 1 bit at a time, sequentially
- ◆ Parallel transmission: 1 word at a time (i.e. n bits in parallel)
- ◆ Advantages of serial transmission: very simple wiring
- ◆ Transmission Characteristics of RS-232
  - ❑ maximum distance of 20 metres
  - ❑ Maximum bit rate 19 600 bps
- ◆ Alternative serial communication standards:
  - ❑ RS-422 (up to 10 Mbps over distance of 1.2km)
  - ❑ USB-1 & USB-2
  - ❑ IEEE 1394 (Firewire / iLink)

# Serial Interfaces: RS-232

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## ◆ Electrical characteristics

- ☐ Logic 1: -3V to -25V; typically -12V
- ☐ Logic 0: +3v to +25V; typically +12V
- ☐ Any signal in the range -3V to +3V has an indeterminate logical state
- ☐ Quiescent or inactive state is -12V (i.e. logic 1)

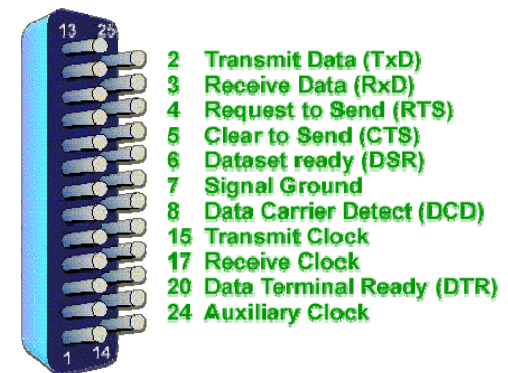
# Serial Interfaces: RS-232

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## ◆ Connectors

- ❑ DB25S is a 25 pin connector with full RS-232 functionality
- ❑ The computer socket has a female outer casing with male connecting pins
- ❑ The terminating cable connector has a male outer casing with female connecting pins

### RS232 Pinout on DB25

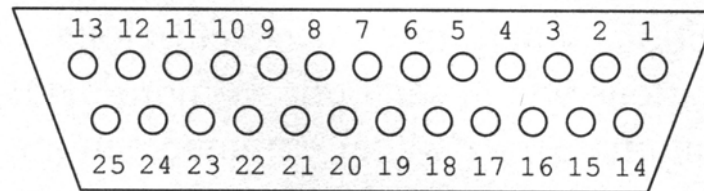


# Serial Interfaces: RS-232

DB 25 MALE DTE	
13CT <sub>s</sub>	25BSY
12Cd	24CLK
11	23DRD
10DC	22RI
9DC	21SQD
8CD	20DTR
7SG	19RT <sub>s</sub>
6DSF	18CLK
5CT <sub>s</sub>	17CLK
4RT <sub>s</sub>	16RXd
3RXD	15CLK
2TXD	14TXd
1SG	
	1 Protective ground
	2 Transmit Data
	3 Receive Data
	4 Request to Send
	5 Clear to Send
	6 Data Set Ready
	7 Signal Ground
	8 Carrier Detect
	9 Positive DC Test
	10 Negative DC Test
	11
	12 Secondary Carrier Detect
	13 Secondary Clear to Send
	14 Secondary Transmit Data
	15 Transmit Clock DCE
	16 Secondary Receive Data
	17 Receiver Clock
	18 Receiver dibit clock
	19 Secondary Request to Snd
	20 Data Terminal Ready
	21 Signal Quality Detector
	22 Ring Indicator
	23 Data Signal Rate Select
	24 Transmit Clock
	25 Busy

# Serial Interfaces: RS-232

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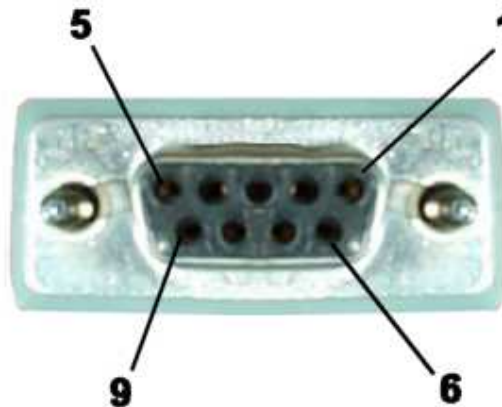
<i>Pin</i>	<i>Signal</i>
2	Tx
3	Rx
4	RTS
5	CTS
6	DSR
7	GND
20	DTR

# Serial Interfaces: RS-232

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## ◆ Connectors

- ❑ DB9S is a 9 pin connector with reduced RS-232 functionality
- ❑ The computer socket has a female outer casing with male connecting pins
- ❑ The terminating cable connector has a male outer casing with female connecting pins



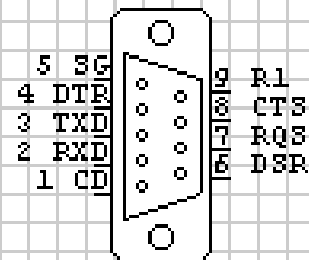


# Serial Interfaces: RS-232

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## RS-232 SERIAL COMMUNICATIONS

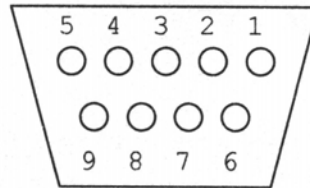
DB9 MALE DTE



1	CD	(in)
2	RXD	(in)
3	TXD	(out)
4	DTR	(out)
5	SG	
6	DSR	(in)
7	RQS	(out)
8	CTS	(in)
9	R1	

# Serial Interfaces: RS-232

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<i>Pin</i>	<i>Signal</i>
2	Rx
3	Tx
4	DTR
5	GND
6	DSR
7	RTS
8	CTS

# Serial Interfaces: RS-232

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## ◆ Connectors

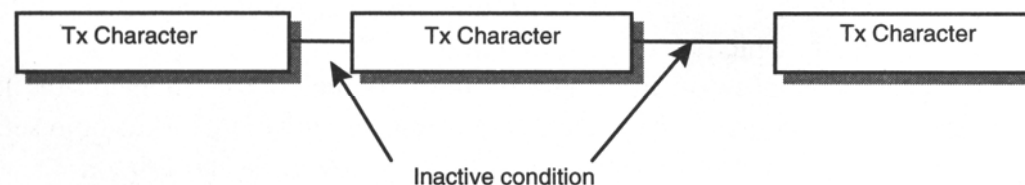
- ☐ Most PCs use either a 9-pin connector for the primary (COM1:) serial port and a 25-pin for a secondary serial port (COM2:), or they use two 9-pin connectors
- ☐ Note: the 25-pin parallel port (LPT1:) is a 25-pin female connector on the PC and male on the cable
- ☐ The serial connector is male on the PC and female on the cable
- ☐ 25-to-9 pin adaptors are available

# Serial Interfaces: RS-232

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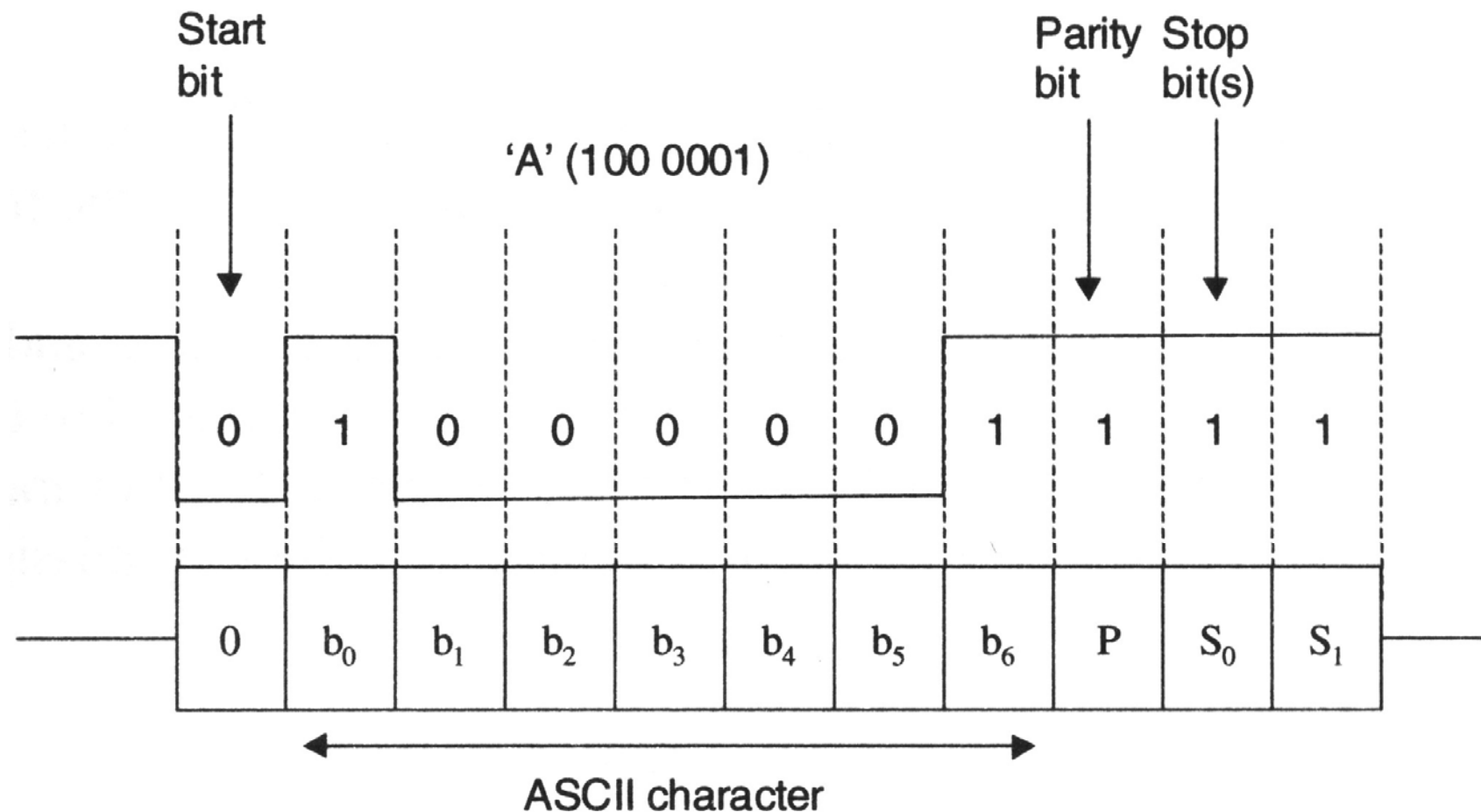
## ◆ Frame format

- ❑ RS-232 uses asynchronous communications
- ❑ Start-stop format



- ❑ Each character is transmitted one at a time
  - Delay between each character: inactive time
  - Inactive time: -12v logic level high
  - Each character is frames by a start bit (0), 7 data bits, 1 parity bit, 2 stop bits
  - ASCII coding is used
  - Parity is either even, odd, none
    - o Parity bit is set to make the total number of either even or odd (or not checked)

# Serial Interfaces: RS-232



# Serial Interfaces: RS-232

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## ◆ Frame format

- ☐ Both the transmitter and the receiver need to be set to the same bit-time interval (baud rate)
- ☐ Since RS-232 is asynchronous, clock rates don't have to be exactly synchronized
- ☐ There is an overhead in using asynchronous communication: the additional start and stop bits
- ☐ The advantages is that it makes communication very simple

# Serial Interfaces: RS-232

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## ◆ Example

ASCII coding, even parity, 2 stop bits:

```
111110100000101100000111111111111100000111111  
11000110011110101001111111111111
```

```
{inactive}11111 {start bit} 0 {'A'}1000001 {parity bit} 0  
{stop bits} 11 {start bit}0 {'p'}0000111 {parity bit} 1  
{stop bits}11 {inactive}11111111 {start bit}0  
{'p'}0000111 {parity bit} 1 {stop bits}11 {inactive}11  
{start bit}0 {'L'}0011001 {parity bit} 1 {stop bits}11
```

Message is 'AppL'

# Serial Interfaces: RS-232

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## ◆ Parity

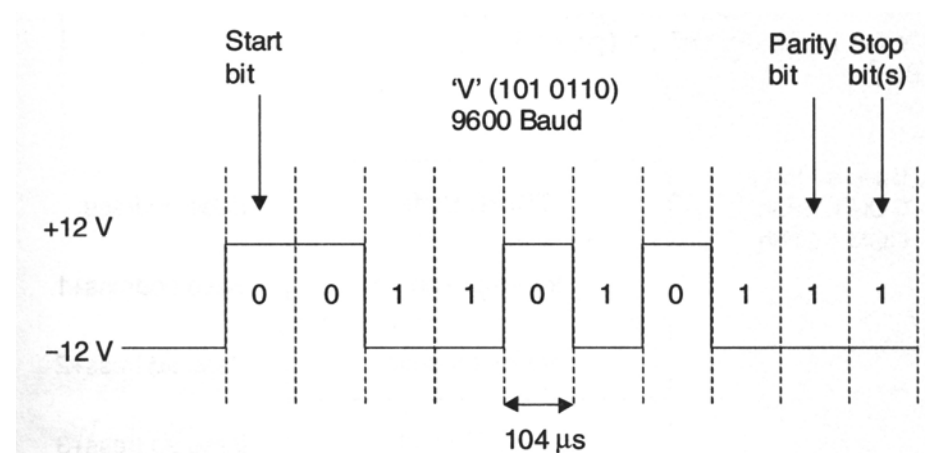
- ☐ Simple form of error coding
- ☐ A parity bit is added to transmitted data to make the number of 1s (in the data) either even (even parity) or odd (odd parity)
- ☐ A single parity bit can only detect an odd number of errors (why?)



# Serial Interfaces: RS-232

## ◆ Baud Rate

- ❑ 11 bits required to send a single character (10 if one stop bits are used)
- ❑ Bit rate (bits/sec): actual rate at which bits are transmitted
- ❑ Baud rate: rate at which the signalling elements, used to represent bits, are transmitted
  - Since one signalling element encodes one bit, both rates are identical
  - However, baud rates will differ from bit rates in modem communication
- ❑ Time period of each bit =  $1/\text{baud rate}$  seconds



# Serial Interfaces: RS-232

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## ◆ Types of equipment

- ❑ DTE Data Terminal Equipment
  - Originally applied to CRT terminals or other input devices
  - Today, DTE mainly applies to a computer
- ❑ DCE Data Communication Equipment
  - Originally applied to modems or similar communications equipment
  - Still applies today
- ❑ A modem is a device that converts a digital signal (e.g. from an RS232 interface) to an analogue signal for transmission over a traditional telephone line (MODEM: MODulator-DEModulator)

# Serial Interfaces: RS-232

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Pin	Name	Abbrev	Functionality
1	Frame Ground	FG	This ground normally connects the outer sheath of the cable to the earth ground
2	Transmit Data	TD	Data is sent from the DTE (computer or terminal) to a DCE via TD
3	Receive Data	RD	Data is send from the DCE to a DTE via RD
4	Request to Send	RTS	DTE sets this active when it is ready to transmit data
5	Clear to Send	CTS	DCE sets this active to inform the DTE that it is ready to receive data
6	Data Set Ready	DSR	Signals that the DCE is ready to communicate
7	Signal Ground	SG	All signals are referenced to the signal ground
20	Data Terminal Ready	DTR	Signals that the DTE is ready to communicate RTS and CTS have no effect is DTR is not asserted and connected to DSR

# Serial Interfaces: RS-232

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## ◆ Communication between two nodes

### □ Handshaking

#### ▪ Hardware handshaking

- RTS
- CTS
- DTR
- DSR

#### ▪ Software handshaking

- Sending special control characters X-OFF, X-On (ctrl-S, ctrl-Q)

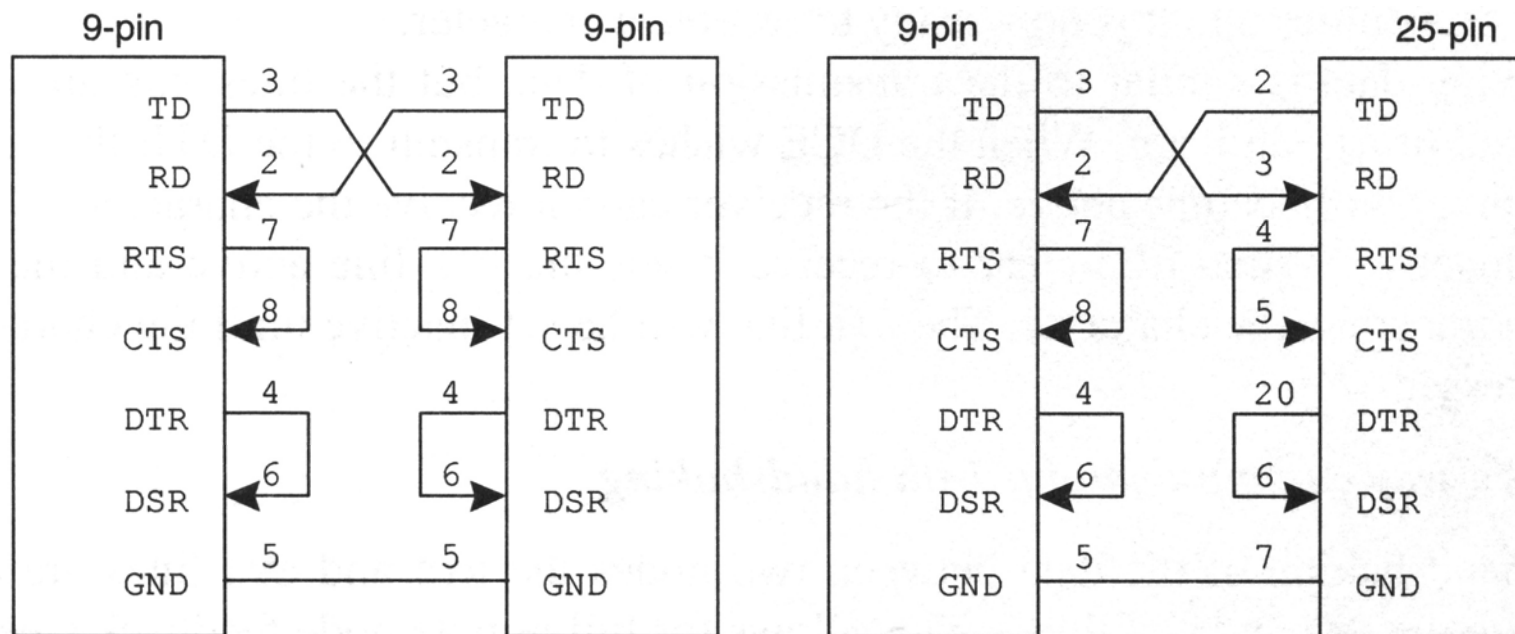
#### ▪ No handshaking

- If no handshaking is used then the receiver must be able to read the received characters before the transmitter send more (otherwise the input buffer will be overwritten)

# Serial Interfaces: RS-232

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- ❑ Simple No Handshaking Connections
  - note the loop-backs



# Serial Interfaces: RS-232

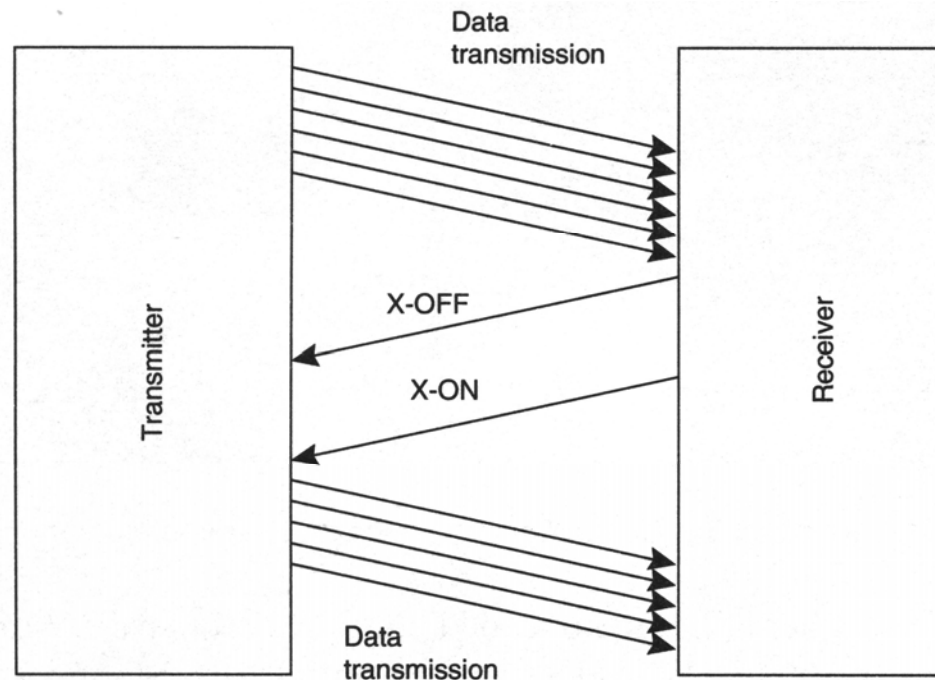
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- ❑ Simple No Handshaking Connections
  - Advantage: no handshaking so fewer wires
  - Disadvantage:
    - o all handshaking has to be done by software (hence, more complex software)
    - o The lack of the DTR means neither device knows whether the other device is powered up and ready for data transfer
  - Typically used when the devices can operate at much faster speeds than the communication channel

# Serial Interfaces: RS-232

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- ❑ Software handshaking  
(Note that the software must recognize the X-ON and X-OFF characters and take the appropriate action)



# Serial Interfaces: RS-232

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- ❑ Hardware handshaking
  - To transmit:
    - o Assert Transmitter RTS (high)
    - o Wait until Transmitter CTS high
    - o When receiver reads from its buffer (buffer empty), Receiver RTS goes high
    - o  $\Rightarrow$  Transmitter CTS, goes high
    - o Transmitter sends data



# Serial Interfaces: RS-232

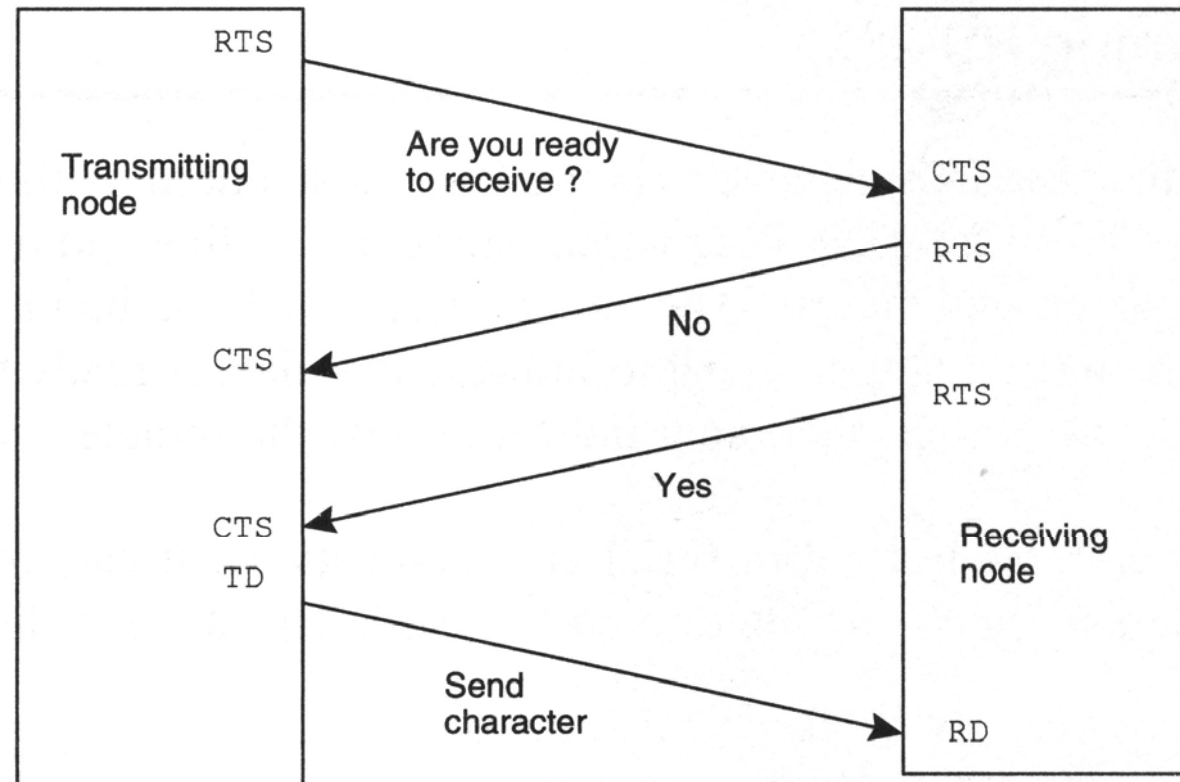
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- ☐ DTR and DSR are responsible for establishing the connection
- ☐ RTS and CTS are responsible for the data transfer
- ☐ Without an active DTR signal, the RTS and CTS signals have no effect
- ☐ Deactivating DTR or DSR breaks the connection

# Serial Interfaces: RS-232

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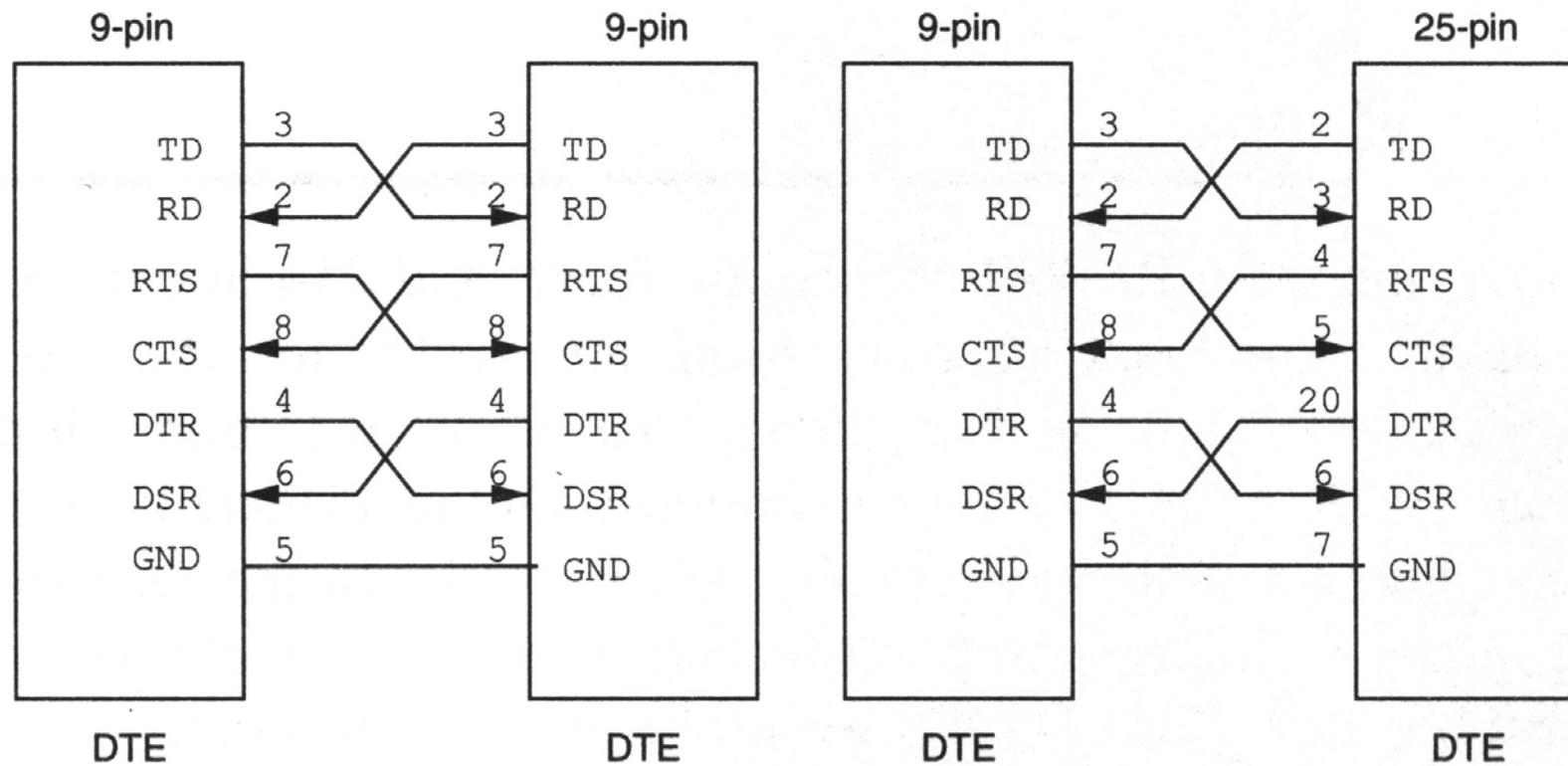
## ❑ Hardware handshaking



# Serial Interfaces: RS-232

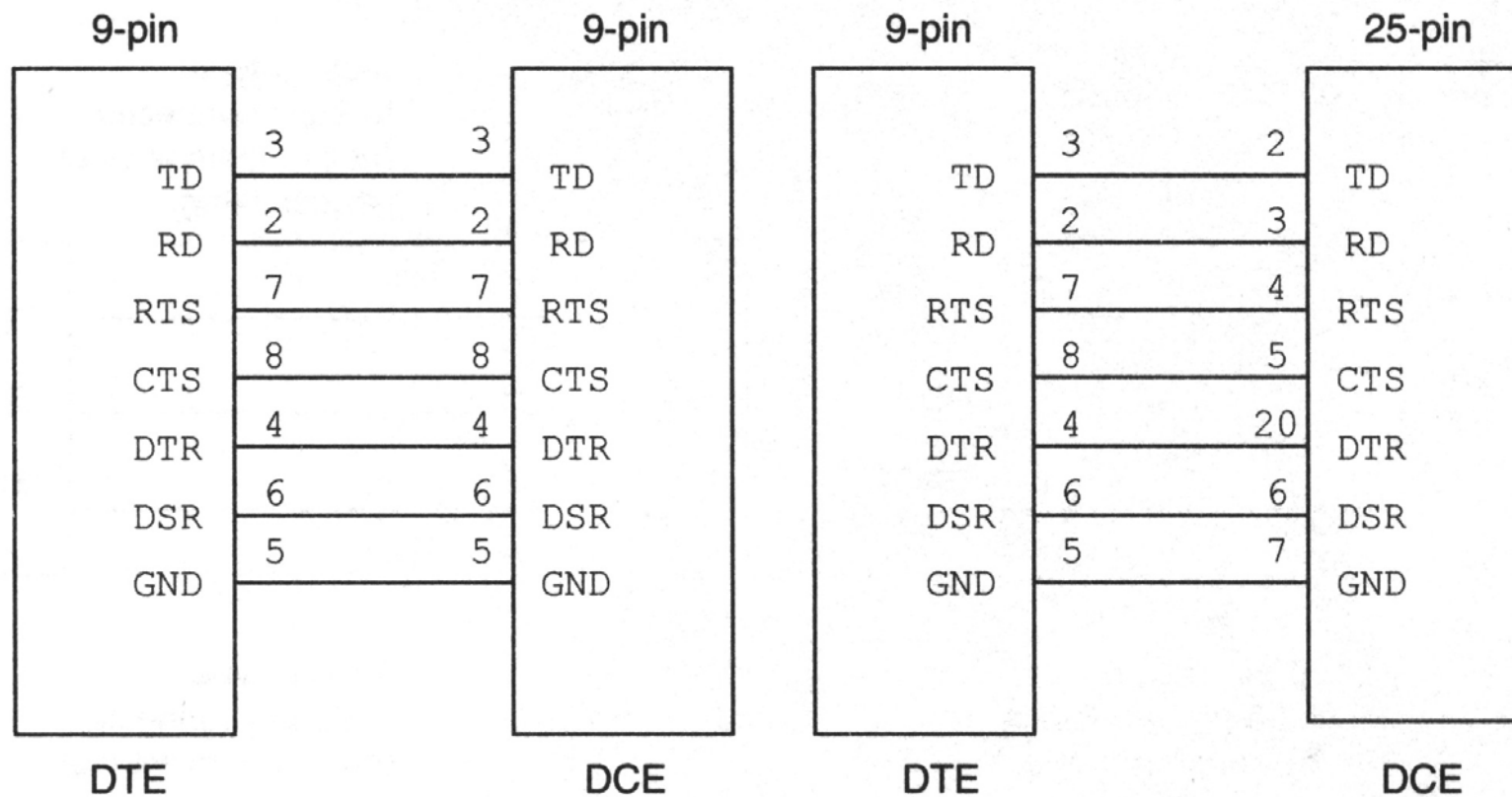
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## ❑ Hardware handshaking



# Serial Interfaces: RS-232

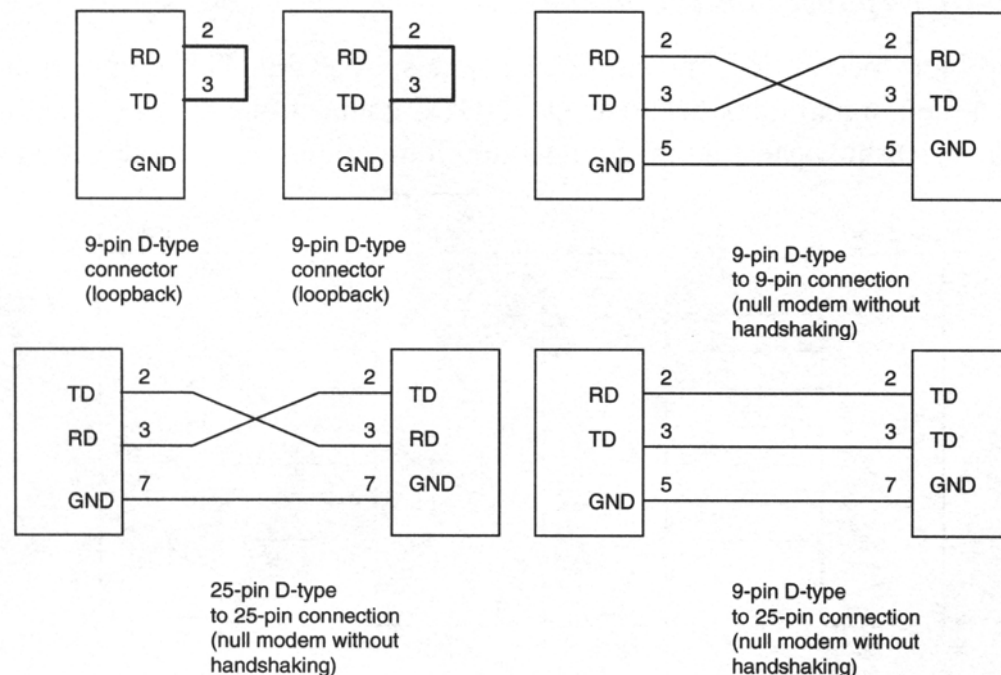
- ❑ DTE to DCE connections for handshaking



# Serial Interfaces: RS-232

## □ Typical System Connections

- Loop-back connections are used to test hardware (cf. program 15.1 in text)
- Null modem connections are used for communication (cf. programs send and receive; 15.2 & 15.3 in text)



# Serial Interfaces: RS-232

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- ❑ Types of error in serial communications
  - Framing error: if the receiver has detected an invalid stop bit then the received serial character does not fit into the frame that the setup data format and the setup baud rate define. Thus the receiver has detected a framing error
  - Break error: if the reception line is at a logical low for a longer time than usual then the receiver assumes that the connection to the transmitter is broken. The transmitter usually drives the line to a logical high level as long as no data is being transferred
  - Overrun error: if data is arriving in the receiver faster than it is read from the receiver buffer register by the CPU, then a later received byte may overwrite the older data not yet read from the buffer.
  - Parity error

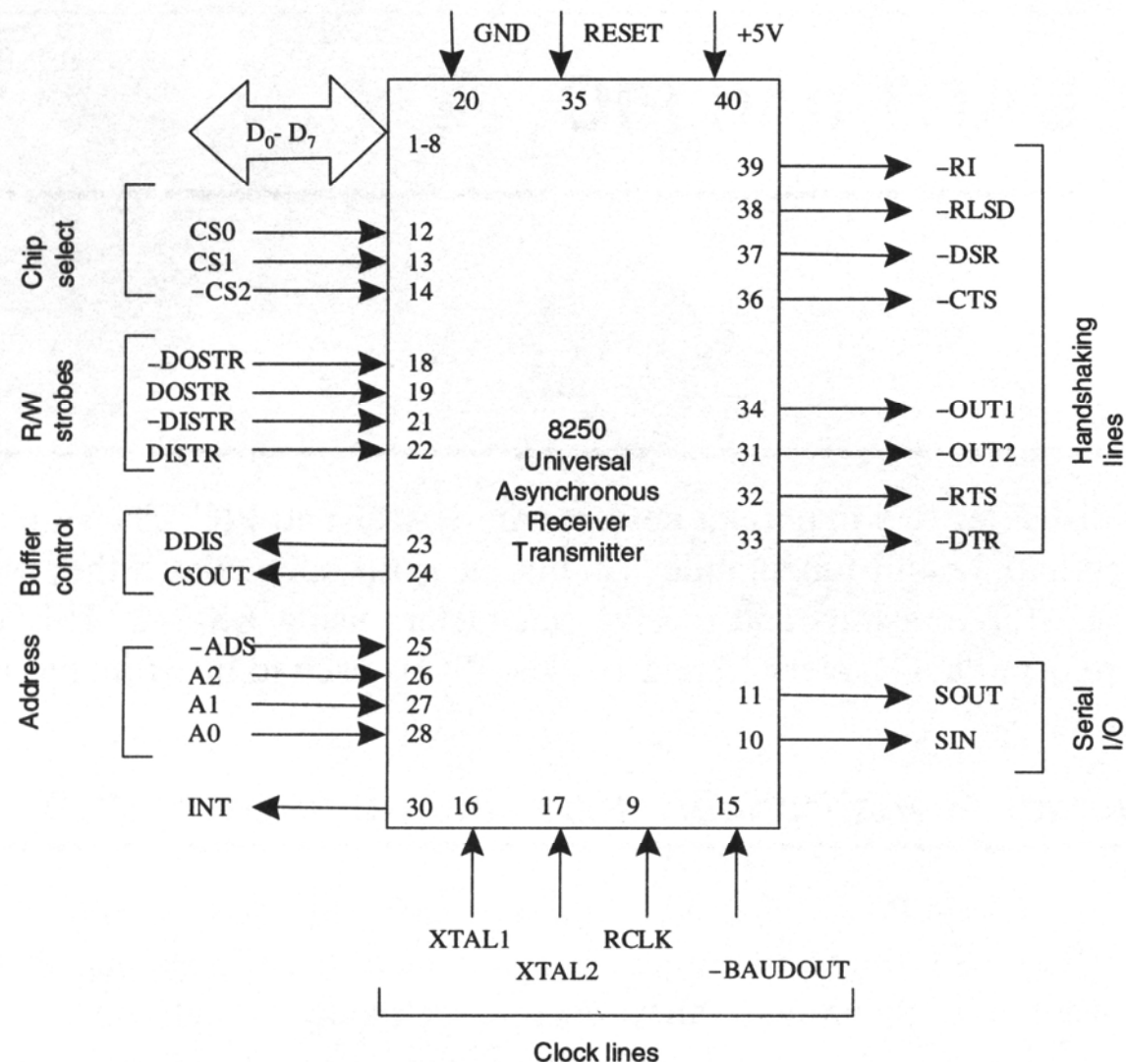
# Serial Devices: UART

---

## ◆ Universal Asynchronous Receiver Transmitter (8250)

- ❑ UART transmits and receives characters using RS-232
- ❑ 40 pin IC
  - Connection to the microprocessor is via D0-D7
  - Microprocessor **write**: asserts DOSTR and  $\overline{\text{DOSTR}}$  (high & low, respectively)
  - Microprocessor **read**: asserts DISTR and  $\overline{\text{DISTR}}$  (high & low, respectively)

# Serial Devices: UART



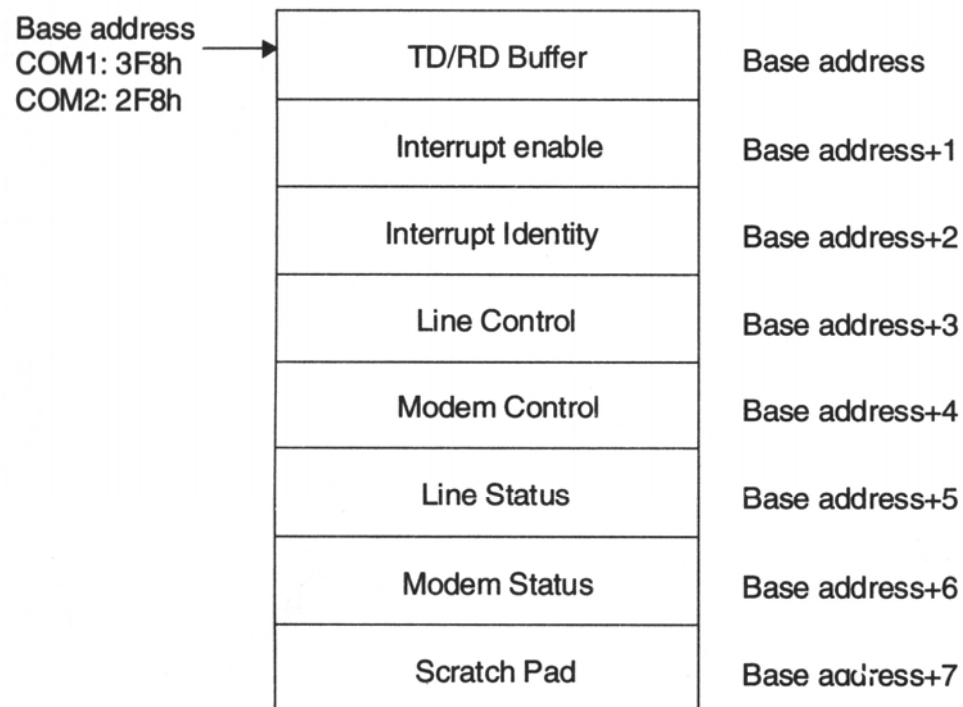


# Serial Devices: UART

---

## ◆ Universal Asynchronous Receiver Transmitter (8250)

- Seven registers, selected using address lines A0, A1, A2



# Serial Devices: UART

---

## ◆ Universal Asynchronous Receiver Transmitter (8250)

- ❑ Main input RS-232 handshaking lines:

$\overline{RI}$	Ring Indicate
$\overline{DSR}$	Data Set Ready
$\overline{CTS}$	Clear To Send

- ❑ Main output RS-232 handshaking lines:

$\overline{RTS}$	Ready to Send
$\overline{DTR}$	Data Terminal Ready

- ❑ Serial Output

SOUT

- ❑ Serial Input

SIN

# Serial Devices: UART

---

## ◆ Universal Asynchronous Receiver Transmitter (8250)

### ☐ Clock input:

XTAL1	connect to a crystal to control the internal
XTAL2	clock oscillator (typically 1.8432 MHz)
BAUDOT	Baud rate (clock frequency / 16N; typically 9600 Baud)

### ☐ Hardware interrupts:

INT

### ☐ Reset

RESET	active low
-------	------------

# Serial Devices: UART

---

## ◆ Programming RS-232

### ❑ Main registers used in RS-232:

- Line Control Register (LCR)
- Line Status Register (LSR)
- Transmit and Receive Buffer (TD/RD)

### ❑ Base Address

- primary port (COM1:) 3F8h
- Secondary port (COM1:) 2F8h
- Standard PC can support up to four COM ports
- The base addresses are set in the BIOS memory; address of each port stored at address locations:

0040 : 0000 (COM1:)

0040 : 0002 (COM2:)

0040 : 0004 (COM3:)

0040 : 0006 (COM4:)

# Serial Devices: UART

---

```
#include <stdio.h>
#include <conio.h>
int main(void){
    int far *ptr; /* 20-bit pointer */
    ptr = (int far *) 0x0400000; /* 0040:0000 */

    clrscr();
    printf("COM1: %04x\n", *ptr);
    printf("COM2: %04x\n", *(ptr+1));
    printf("COM3: %04x\n", *(ptr+2));
    printf("COM4: %04x\n", *(ptr+3));
    return(0);
}
```

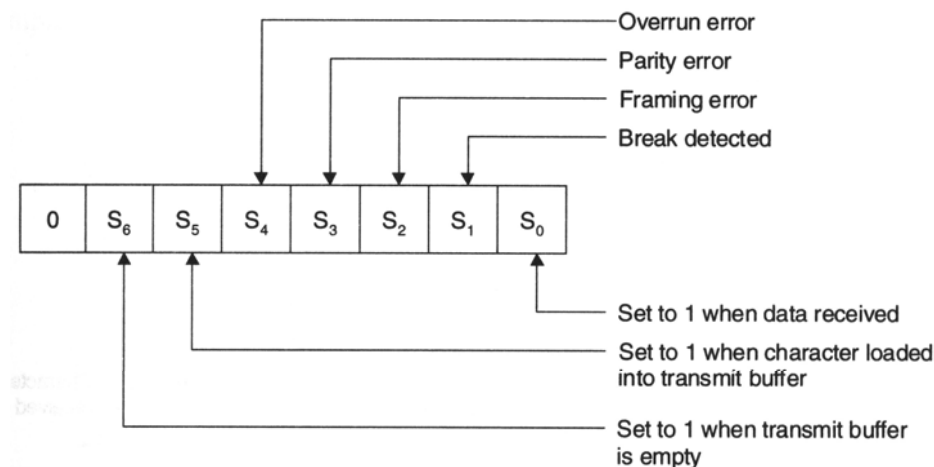
# Serial Devices: UART

---

## ◆ Programming RS-232

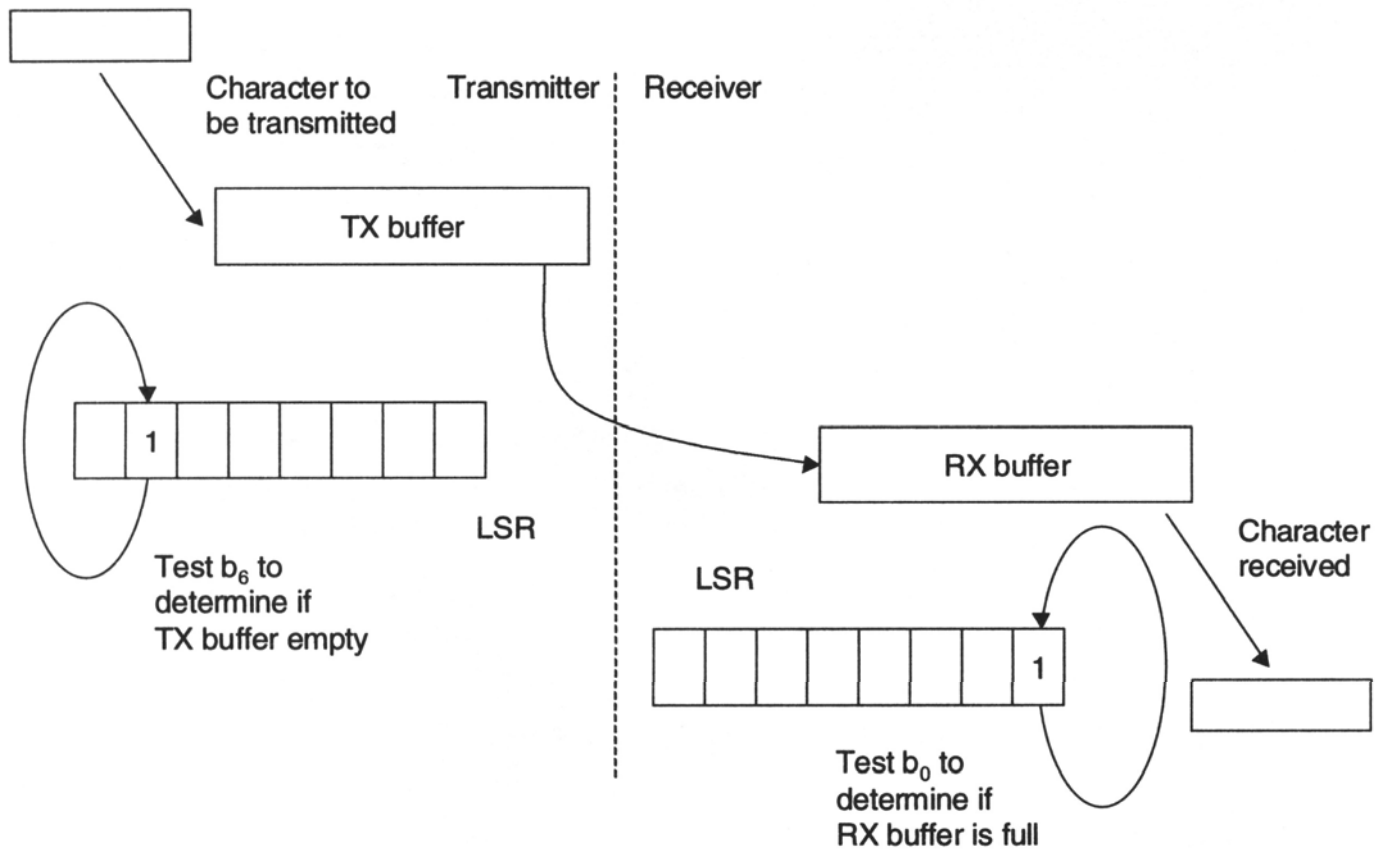
### □ Line Status Register (LSR)

- Determines the status of the transmitter and receiver buffers
- Read-only; all bits set by hardware



- Status bit S<sub>6</sub> should be checked to see if the output buffer (TD/RD) is empty – i.e. data has been sent – before writing to the buffer

# Serial Devices: UART



# Serial Devices: UART

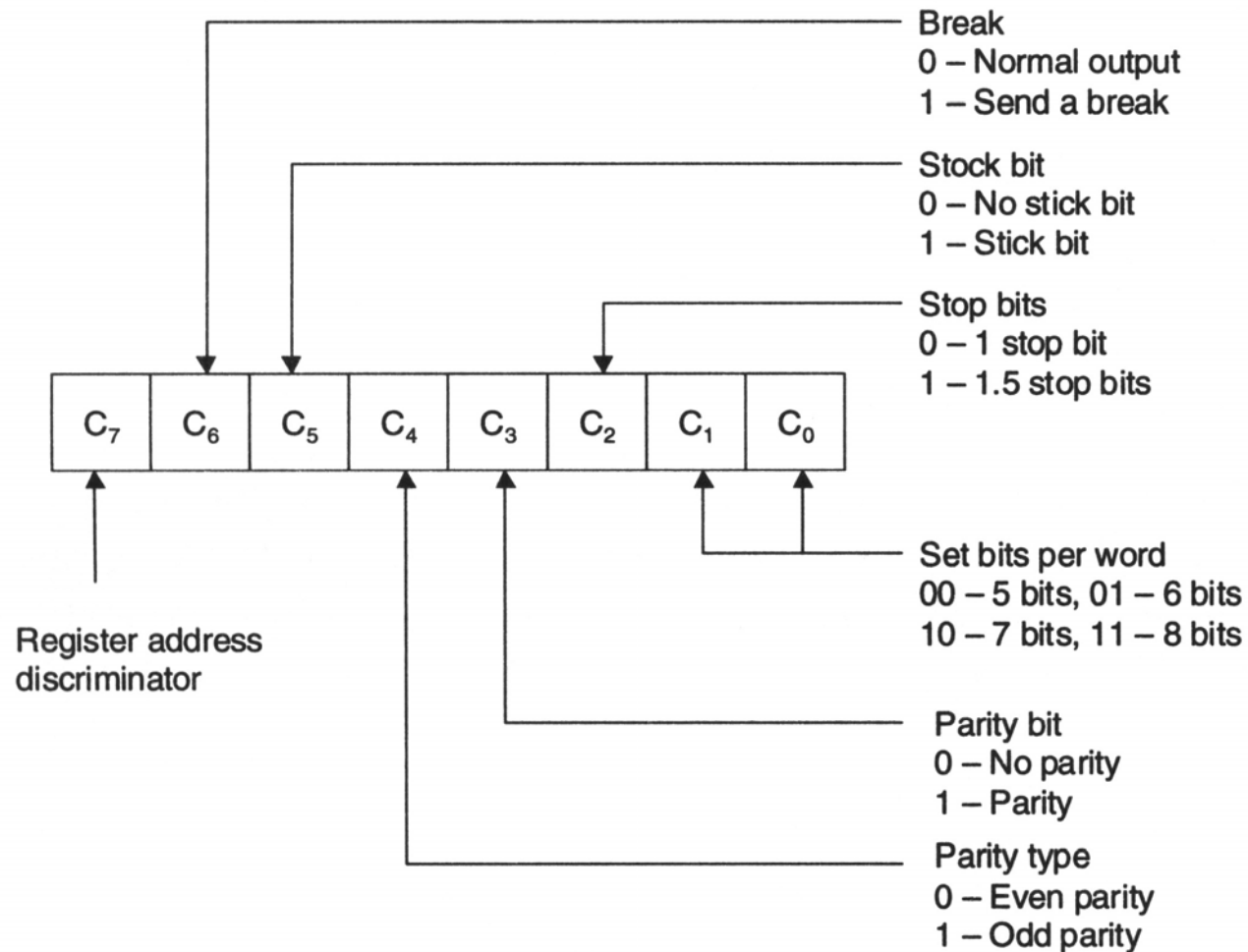
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## ◆ Programming RS-232

- Line Control Register (LCR)
  - Read/Write register
  - Sets up communications parameters
    - o Number of bits per character
    - o Parity
    - o Number of stop bits
    - o Baud rate



# Serial Devices: UART



# Serial Devices - UART

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## ◆ Programming RS-232

### □ Line Control Register (LCR)

#### ▪ Baud rate

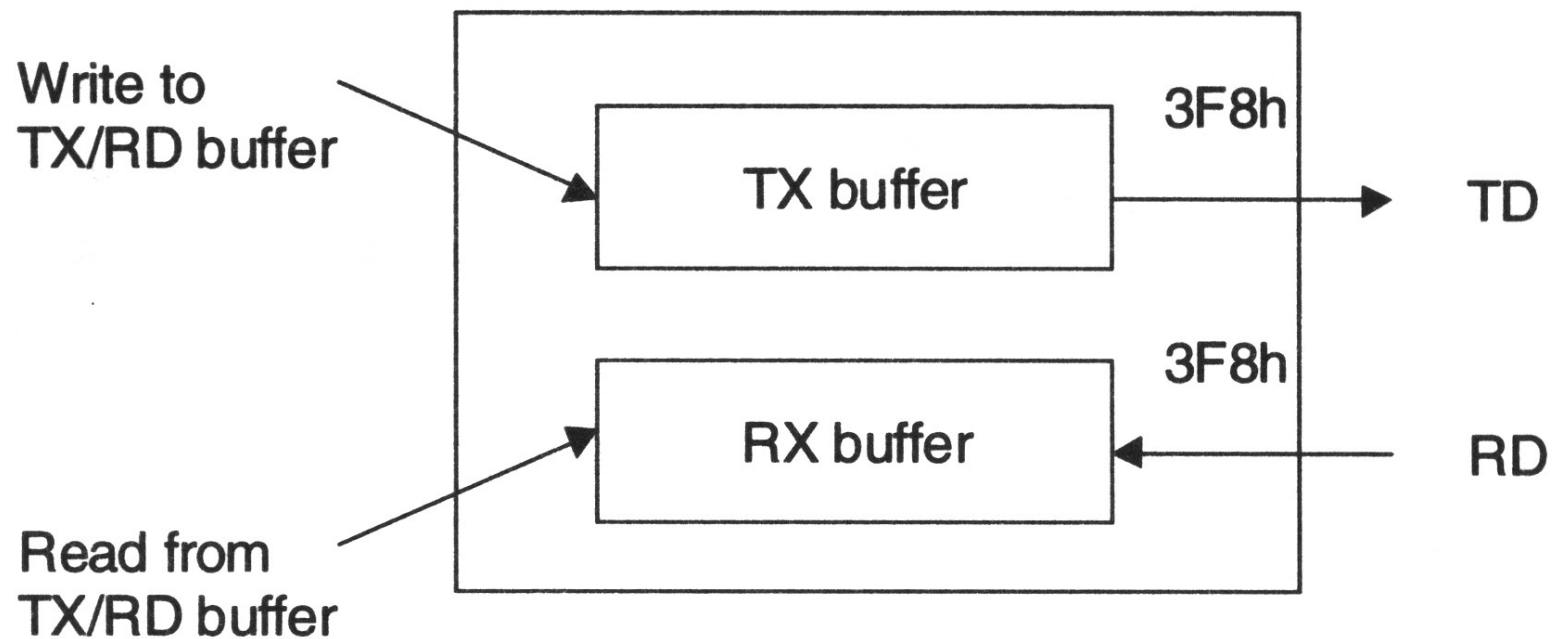
- C7 set high to access the Baud rate divider
- (C7 set low to access the TR/RX buffer)
- Set baud rate by loading 16-bit divisor N  
LSB (least significant byte) -> TR/RX buffer address  
MSB (most significant byte)-> TR/RX buffer address + 1
- Value loaded depends on crystal frequency connected to the IC

$$\text{Baud rate} = \text{Clock frequency} / (16 * N)$$

e.g., for 1.8432 MHz crystal frequency, 9600 baud,  
 $N = 1.8432\text{E}6 / (9600 * 16) = 12 \text{ (000Ch)}$

# Serial Devices: UART

---



# Serial Interfaces: RS-232

---

```
/*      send.c                                          */

#define  TXDATA          0x3F8
#define  LSR             0x3FD
#define  LCR             0x3FB

#include<stdio.h>
#include <conio.h> /* included for getch              */
#include <dos.h>   /* included for inputb and outputb  */

void setup_serial(void);
void send_character(int ch);

int  main(void)
{
    int ch;
    puts("Transmitter program. Please enter text (Cntl-D to end)");
    setup_serial();
    do
    {
        ch=getch();
        send_character(ch);
    } while (ch!=4);
    return(0);
}
```

# Serial Interfaces: RS-232

---

```
void setup_serial(void)
{
    /* set up bit 7 to a 1 to set Register address bit */

    outportb( LCR, 0x80);

    /* load TxRegister with 12, crystal frequency is 1.8432MHz */

    outportb(TXDATA,0x0C);
    outportb(TXDATA+1,0x00);

    /* Bit pattern loaded is 00001010b, from msb to lsb these are: */
    /* Access TD/RD buffer, normal output, no stick bit */
    /* even parity, parity on, 1 stop bit, 7 data bits */
    outportb(LCR, 0x0A);
}

void send_character(int ch)
{
    char status;

    /*repeat until Tx buffer empty ie bit 6 set*/
    do {
        status = inportb(LSR) & 0x40;
    } while (status!=0x40);

    outportb(TXDATA,(char) ch);
}
```

# Serial Interfaces: RS-232

---

```
/*    receive.c                                */

#define TXDATA      0x3F8
#define LSR         0x3FD
#define LCR         0x3FB
#include <stdio.h>
#include <conio.h>    /* included for getch          */
#include <dos.h>      /* included for inputb and outputb */

void      setup_serial(void);
int       get_character(void);

int main(void)
{
    int inchar;
    setup_serial();
    do
    {
        inchar=get_character();
        putchar(inchar);
    } while (inchar!=4);
    return(0);
}
```

# Serial Interfaces: RS-232

---

```
void setup_serial(void)
{
    /* set up bit 7 to a 1 to set Register address bit */

    outportb( LCR, 0x80);

    /* load TxRegister with 12, crystal frequency is 1.8432MHz */

    outportb(TXDATA,0x0C);
    outportb(TXDATA+1,0x00);

    /*      Bit pattern loaded is 00001010b, from msb to lsb these are: */
    /*      Access TD/RD buffer, normal output, no stick bit */
    /*      even parity, parity on, 1 stop bit, 7 data bits */
    outportb(LCR, 0x0A);
}

int get_character(void)
{
    int status;
    /* Repeat until bit 1 in LSR is set */
    do {
        status = inportb(LSR) & 0x01;
    } while (status!=0x01);

    return( (int)inportb(TXDATA));
}
```





# Serial Interfaces: RS-422

---

- ◆ Signal levels 0 to +5V
- ◆ Uses differential amplifiers at both transmitting and receiving ends to achieve high noise immunity
- ◆ Up to 10 Mbps over distance of 1.2km





# Universal Serial Bus - USB 2.0

---

## ◆ Motivation for the development of the USB

- ❑ Cost, configuration, and attachment of many peripheral devices to a PC creates problems
- ❑ Traditional interfacing (legacy systems)
  - Peripherals mapped onto the I/O space
    - o Although the I/O space has 64k of addressable locations (16 bits), legacy ISA bus PCs only decode 10 bit, leaving only 1k of I/O space
    - o I/O address conflicts are common
  - Assigned a specific IRQ (Interrupt Request line)
    - o But there is a limited number of IRQs available (16)
    - o Some of these are dedicated to specific devices
    - o Some are shared (but not at the same time)
  - Sometimes assigned a DMA (Direct Memory Access) channel

# Universal Serial Bus - USB 2.0

---

## ◆ Motivation for the development of the USB

### □ Traditional interfacing (legacy systems)

#### ▪ Physical limitation:

- o serial and parallel interfaces only support a single device each
- o Only solution is to add additional (expensive) expansion cards

#### ▪ Usability

- o Too many connectors and/or cables
- o System must be shut down to attach most peripheral
- o System must be restarted to install/load software

# Universal Serial Bus - USB 2.0

---

## ◆ USB Design Goals

- ☐ Single connector type for any peripheral
- ☐ Ability to attach many peripherals
- ☐ Ease system resource conflicts
- ☐ Hot plug support (hot-swappable)
- ☐ Automatic detection and configuration (plug-and-play)
  - Dynamically loadable and unloadable drivers
    - o Plug in device
    - o Host detects device
    - o Host interrogates device
    - o Host loads appropriate driver
- ☐ Low-cost for both PC and peripheral
- ☐ Enhanced performance
- ☐ Support for new peripheral designs
- ☐ Support for legacy hardware and software
- ☐ Low-power implementation

# Universal Serial Bus - USB 2.0

---

- ◆ 650 pages in the definition of the USB 2.0 specification (see [www.usb.org](http://www.usb.org))
- ◆ USB 1.1 supported two speeds:
  - ☐ Full speed 12 Mbit/s
  - ☐ Low speed 1.5 Mbits/s
    - Less susceptible to EMI (electromagnetic interference)
    - Can use cheaper components
- ◆ USB 2.0 supports 480 Mbits/s
  - ☐ High speed
  - ☐ Competes with IEEE 1394 (Firewire) serial bus

# Universal Serial Bus - USB 2.0

---

## ◆ USB is host controlled

- ☐ Only one host per bus
- ☐ Host is responsible for
  - undertaking all transactions
  - Scheduling bandwidth
  - Data can be sent by one of several transaction methods using a token-based protocol

## ◆ USB uses a tiered star topology

- ☐ Devices/peripherals can't be daisy-chained
- ☐ Need a hub
  - Some peripherals have in-built hubs (e.g. keyboards)
  - You can plug other devices into these
- ☐ Up to 127 devices can be connected to any one USB at any one given time



# Universal Serial Bus - USB 2.0

---

## ◆ USB host controller specifications (USB 1.1)

### ☐ UHCI (Universal Host Controller Interface)

- Developed by Intel
- More work is done in software
- Hence cheaper hardware

### ☐ OHCI (Open Host Controller Interface)

- Developed by Compaq, Microsoft, and National Semiconductor
- More work done in hardware
- Hence simpler software!

### ☐ USB 2.0

- EHCI (Enhanced Host Controller Interface)

# Universal Serial Bus - USB 2.0

---

## ◆ USB Transfer modes

- ☐ Control
- ☐ Interrupt
- ☐ Bulk
- ☐ Isochronous
  - Allows a device to reserve a defined amount of bandwidth with guaranteed latency
  - Ideal for audio or video
  - Avoids dropped frames due to congestion

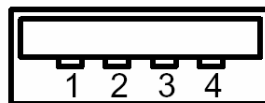
# Universal Serial Bus - USB 2.0

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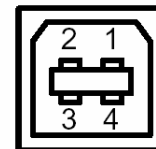
## ◆ Connectors

- ☐ Upstream connection to the host
- ☐ Downstream connection to a device
- ☐ Not interchangeable
- ☐ Type A plugs (and sockets) always face upstream to the host
- ☐ Type B plugs (and sockets) always face downstream to the peripheral device
- ☐ USB 2.0 includes mini-USB B connectors

Receptacle Type A



Receptical Type B



# Universal Serial Bus - USB 2.0

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## ◆ Connectors

Pin Number	Cable Colour	Function
1	Red	V <sub>BUS</sub> (5 volts)
2	White	D-
3	Green	D+
4	Black	Ground

# Universal Serial Bus - USB 2.0

---

## ◆ Electrical Characteristics

- ☐ USB uses a differential transmission pair for data
- ☐ Differential '1' is transmitted by
  - Pulling D+ over 2.8V
  - Pulling D- under 0.3V
- ☐ Differential '0' is transmitted by
  - Pulling D- over 2.8V
  - Pulling D+ under 0.3V
- ☐ The polarity of the signal is inverted depending on the speed of the bus
  - Use 'J' and 'K' to signify logic levels
  - 'J' is differential 0 in low speed
  - 'J' is differential 1 in high speed

# Universal Serial Bus - USB 2.0

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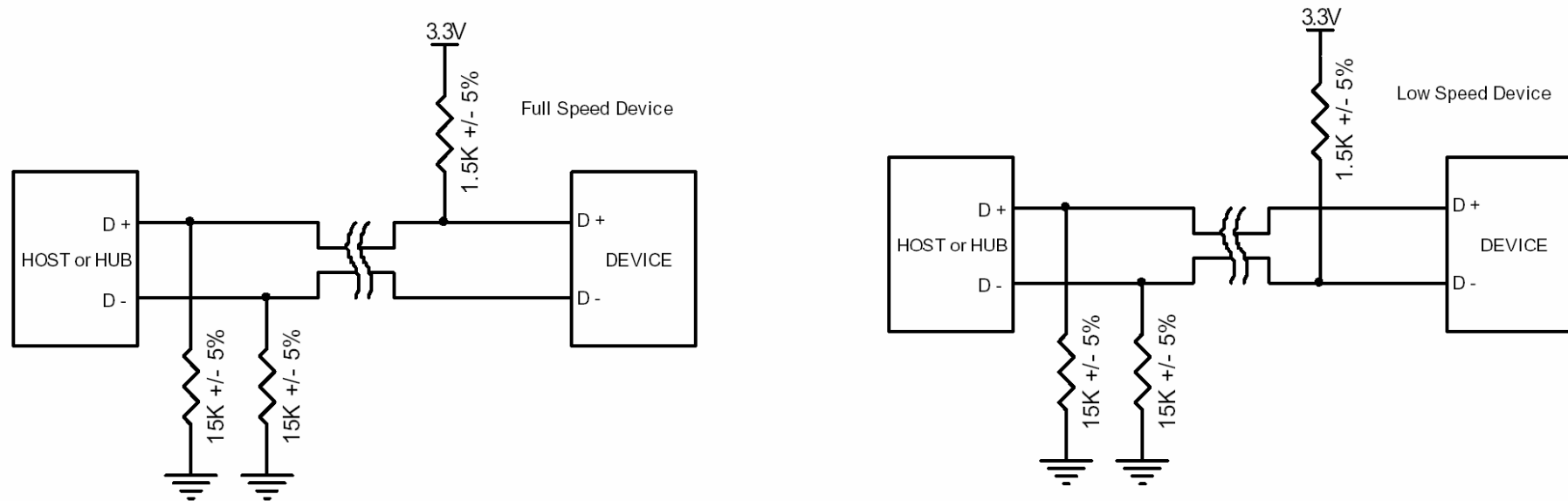
## ◆ Electrical Characteristics

- ☐ USB also uses single ended outputs (on D+, D-, or both)
- ☐ For example:
  - Single Ended Zero (SE0)
  - Used to signify a device reset if held for more than 10mS
  - Generated by holding both D- and D+ low ( $<0.3V$ )
- ☐ If you are building a USB interface / device, you can't get away with just sampling the differential output

# Universal Serial Bus - USB 2.0

## ◆ Speed Identification

- ❑ A USB device must indicate its speed by pulling either the D+ or D- line high to 3.3V (full speed and low speed respectively)



# Universal Serial Bus - USB 2.0

---

## ◆ Speed Identification

- ☐ High speed devices start by connecting as full speed
- ☐ Once attached, it does a high speed chirp (signal with a time-varying increase in frequency) during reset
- ☐ A high speed connection is then established if the hub supports it

## ◆ Presence of a device on the bus

- ☐ These pull-up resistors are also used by the host or hub to detect the presence of a device connected to its port
- ☐ Without a pull-up resistor, USB assumes there is nothing connected to the bus
- ☐ Some devices have the resistor built into its silicon, which can be turned on and off under firmware control; others require an external resistor



# Universal Serial Bus - USB 2.0

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## ◆ Speed Identification

- ❑ A USB 2.0 compliant device (downstream) is not required to support high-speed mode
  - Allows cheaper devices to be produced if speed isn't critical
  - Same is true for low speed USB 1.1 devices: don't have to support full speed
- ❑ High speed device MUST NOT support low speed mode
  - Only support full speed (to establish connection) and high speed (if successfully negotiated with the host)
- ❑ A USB 2.0 compliant downstream facing device (hub or host) must support all three modes
  - High speed 480 Mbits/s
  - Full speed 12 Mbits/s
  - Low speed 1.5 Mbits/s

# Universal Serial Bus - USB 2.0

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## ◆ Power Characteristics

- ❑ One big advantage of USB is that devices can be powered by the bus (bus-powered devices)
- ❑ But there are restrictions
  - USB devices specify power consumption in 2mA units in the configuration descriptor (later)
  - A device cannot increase its power consumption over that specified during enumeration (set-up phase) EVEN IF IT LOSES EXTERNAL POWER
  - 3 classes of USB power functions
    - o Low-power bus powered functions
    - o High-power bus powered functions
    - o Self-powered functions

# Universal Serial Bus - USB 2.0

---

## ◆ Power Characteristics

- ❑ Low-power bus powered functions
  - Devices draw all its power from the  $V_{BUS}$
  - Cannot draw more than one unit load (100mA)
  - Must be designed to work with  $4.40V \leq V_{BUS} \leq 5.25V$
- ❑ High-power bus powered functions
  - Device draws all its power from the  $V_{BUS}$
  - Cannot draw more than one unit load (100mA) until it has been configured
  - Then it can draw 5 unit loads maximum, provided it requested it in its descriptor
- ❑ Self-powered functions
  - May draw up to 1 unit load from the bus and derive the rest from an external source
  - The 1 unit load allows detection and enumeration without external supply (but not operation)

# Universal Serial Bus - USB 2.0

---

## ◆ Suspend Mode

- ❑ Suspend mode is mandatory on all devices
- ❑ Suspend mode imposes further constraints on power
- ❑ The maximum suspend current is proportional to the unit load
  - 1 unit load device: maximum suspend current is 500 $\mu$ A
  - This is not a lot!
  - The pull-up resistor (at the device, for detection) and pull-down resistor (at the host) will draw 200 $\mu$ A ( $3.3\text{V} / 16.5\text{k}\Omega$ )
  - 5V to 3.3V regulators to allow 3.3V devices to work will draw more

# Universal Serial Bus - USB 2.0

---

## ◆ Suspend Mode

- ☐ A USB device will enter suspend when there is no activity on the bus for greater than 3.0ms
- ☐ It then has a further 7ms to shutdown the device and draw no more than the suspend current
- ☐ To stay connected, the device must still provide power to the speed selection pull-up resistor during suspend
- ☐ USB can send 'start of frame' packets sent periodically to prevent an idle (no data) bus entering suspend
  - High speed bus: frames sent every 124  $\mu$ s
  - Full speed bus: frames sent every 1ms
  - Low speed bus: EOP (end of packet) every 1 ms
- ☐ A suspended device will resume operation when it receives any non-idle signalling.

# Universal Serial Bus - USB 2.0

---

## ◆ USB Protocols - Overview

- ❑ RS-232 serial interface does not define the format of the data being sent
- ❑ USB is made up of several layers of protocols
  - ICs normally take care of the lower layer
- ❑ Each USB transaction consists of
  - Token Packet (header defining what will follow)
  - Optional Data Packet
  - Status Packet (to acknowledge transaction & provide error correction)
- ❑ Remember: USB is host-centric – the host initiates all transactions

# Universal Serial Bus - USB 2.0

---

## ◆ USB Protocols - Overview

- ❑ Host generates token packet:
  - Specifies what is to follow, and
  - Whether the data transaction is a read or a write
- ❑ Host (usually) then sends a data packet
- ❑ Followed by a handshaking packet (check for errors)
- ❑ *Note: data is transmitted LSB first*

# Universal Serial Bus - USB 2.0

---

## ◆ USB Packet Structure

- Each USB packet comprises the following fields
  - Sync
    - 8 bits long
    - used to synchronize the clock of the receiver with the clock of the transmitter
  - PID (Packet ID)
    - identifies the type of packet being sent
    - 4 bits (but complemented and repeated to give 8 bits)
  - ADDR
    - Address of device to which packet is being sent
    - 7 bits (127 devices)
    - Address 0 is not valid
      - any device which is not yet assigned an address must respond to packets sent to address zero



# Universal Serial Bus - USB 2.0

---

## ◆ USB Packet Structure

- Each USB packet comprises the following fields
  - ENDP
    - End point field
    - 4 bits (16 end-points)
  - CRC
    - Cyclic Redundancy Check
    - Performed on all the data within the packet payload
    - Token packets have a 5 bit CRC
    - Data packets have a 16 bit CRC
  - EOP
    - End of packet
    - Signalled by a Single Ended Zero (SE0) for approx. 2 bit times followed by a 'J' for 1 bit time

# Universal Serial Bus - USB 2.0

---

## USB Packet Structure – PID Values

Group	PID Value	Packet Identifier
Token	0001	OUT Token
	1001	IN Token
	0101	SOF Token
	1101	SETUP Token
Data	0011	DATA0
	1011	DATA1
	0111	DATA2
	1111	MDATA
Handshake	0010	ACK Handshake
	1010	NAK Handshake
	1110	STALL Handshake
	0110	NYET (No Response Yet)
Special	1100	PREAmble
	1100	ERR
	1000	Split
	0100	Ping

PID <sub>0</sub>	PID <sub>1</sub>	PID <sub>2</sub>	PID <sub>3</sub>	nPID <sub>0</sub>	nPID <sub>1</sub>	nPID <sub>2</sub>	nPID <sub>3</sub>
------------------	------------------	------------------	------------------	-------------------	-------------------	-------------------	-------------------

# Universal Serial Bus - USB 2.0

---

## ◆ USB Packet Types

- USB has four packet types
  - Token packets (type of transaction)
  - Data Packets (payload / information)
  - Handshake Packets (ack & error correction)
  - Start of Frame packets (flag start of a new frame)

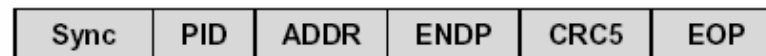
# Universal Serial Bus - USB 2.0

---

## ◆ USB Packet Types

### □ Token packets

- 3 types of token packets
- **In** informs the USB device that host wishes to read info.
- **Out** Informs the USB device that host wishes to send info.
- **Setup** used to begin control transfers
- Packet format:



# Universal Serial Bus - USB 2.0

---

## ◆ USB Packet Types

- Data packets
  - 2 types of data packets
  - Each can transmit 0-1023 bytes of data
  - **Data0**
  - **Data1**
  - Packet format:



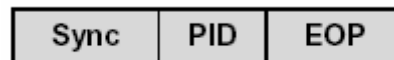
# Universal Serial Bus - USB 2.0

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## ◆ USB Packet Types

### ☐ Handshake packets

- 3 types of handshake packets
- **ACK** Acknowledgement that the packet has been successfully received
- **NAK** Reports that the device can neither send nor receive data
- **STALL** The device needs intervention from the host
- Packet format:



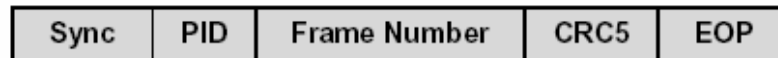
# Universal Serial Bus - USB 2.0

---

## ◆ USB Packet Types

### □ Start of Frame packets

- Consists of an 11-bit number & is sent by the host every 1mS +/- 500nS
- Packet format:



# Universal Serial Bus - USB 2.0

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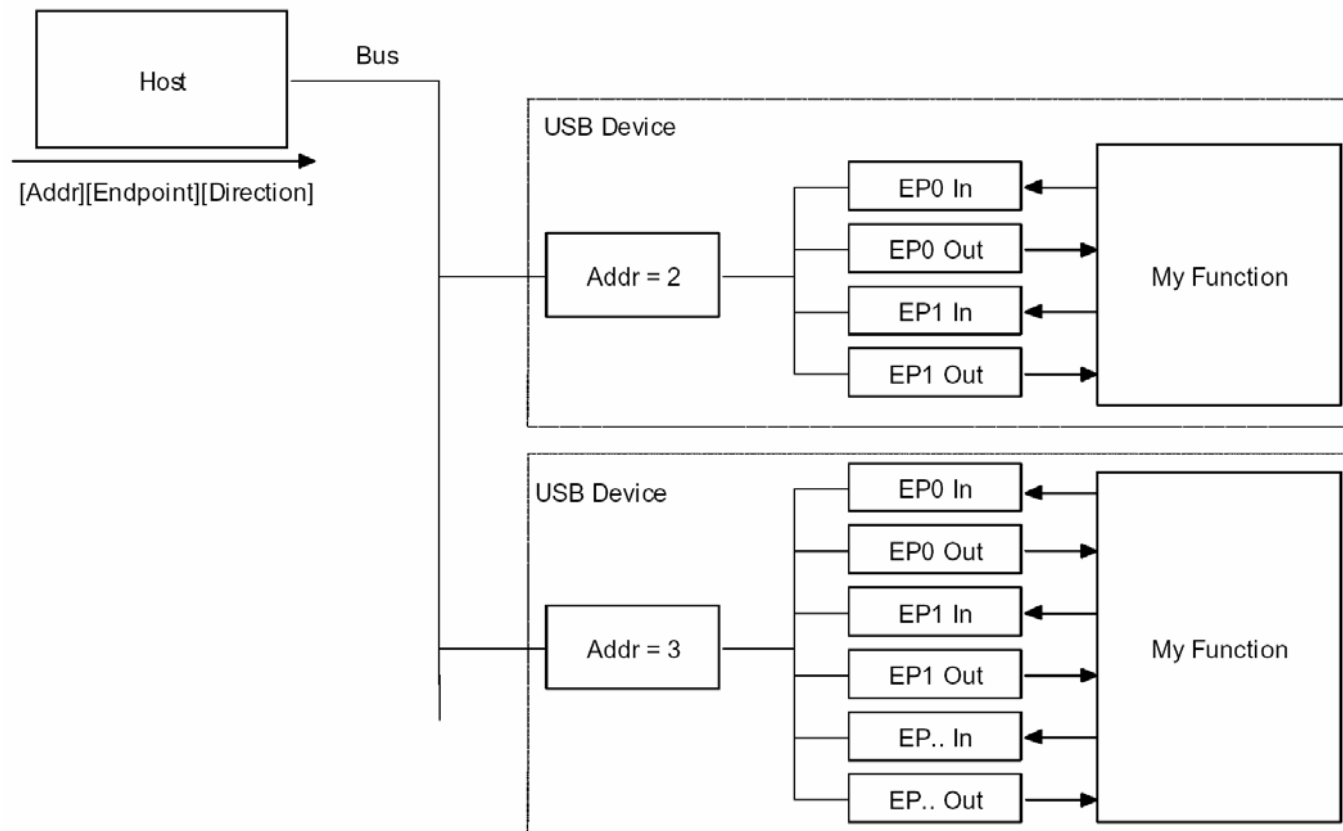
## ◆ USB Functions

- ❑ USB functions are USB devices / peripherals (e.g. printers, disk drives, cameras, scanners, modems, ...)
- ❑ Most functions have a series of buffers, typically 8 bytes long
- ❑ Each buffer will belong to an endpoint
  - **EP0 IN**
  - **EP0 OUT**
  - **EP1 IN**
  - ...



# Universal Serial Bus - USB 2.0

---



# Universal Serial Bus - USB 2.0

---

## ◆ USB Functions – Conceptual Example

- ☐ Host sends a device descriptor request
- ☐ Function hardware reads the setup packet
- ☐ Function determines from the address field whether the packet is addressed to it
- ☐ If it is, the function copies the payload of the following data packet to the appropriate endpoint buffer (given by the value in the endpoint field of the setup token)
- ☐ Function then sends a handshake packet to acknowledge the reception of the byte
- ☐ Function generates an internal interrupt for the appropriate endpoint, flagging that it has received a packet
- ☐ Typically, this is all done in hardware
- ☐ The software interrupt handler now reads the endpoint buffer, interprets it, and initiates the appropriate action

# Universal Serial Bus - USB 2.0

---

## ◆ Endpoints

- ☐ Sources or sinks of data
- ☐ Occur at the end of the communications channel at the USB function
- ☐ The endpoint OUT is where the host sends data to (on the function). E.g. **EP1 OUT**
- ☐ The endpoint IN is where the function places data when it is to be sent to the host, e.g. **EP1 IN**
- ☐ NB: the function cannot simply write to the bus – the bus is controlled by the host and it must send an **IN** packet to enable the transfer of data from **EP2 IN** to the host

# Universal Serial Bus - USB 2.0

---

## ◆ Pipes

- ❑ A pipe is a logical connection between the host and the endpoint(s)
- ❑ Client software transfers data through pipes
- ❑ Pipes have a set of parameters associated with them
  - How much bandwidth is allocated
  - What transfer type it uses
    - o Control
    - o Bulk
    - o Isochronous
    - o Interrupt
  - Direction of data flow
  - Maximum packet/buffer size
- ❑ Default pipe is bidirectional, between **EP0 IN** and **EP0 OUT** with control transfer type

# Universal Serial Bus - USB 2.0

---

## ◆ Pipes

- ❑ USB defines two types of pipes
- ❑ Stream Pipes
  - No defined format
  - Can send any type of data down a stream pipe
  - Data flows sequentially
  - Data has a predefined direction (in or out)
  - Support bulk, isochronous, interrupt transfer types
  - Can be controlled by either host or function / device
- ❑ Message Pipes
  - Defined format
  - Host controlled
  - Initiated by a request sent from the host
  - Data is then sent in the desired direction (depends on the request)
  - Supports only control transfers

# Universal Serial Bus - USB 2.0

---

## ◆ Endpoint / Transfer Types

- ☐ Control Transfers
- ☐ Interrupt Transfers
- ☐ Isochronous Transfers
- ☐ Bulk Transfers

# Universal Serial Bus - USB 2.0

---

## ◆ Control Transfers

- ☐ Used for command and status operations
- ☐ E.g. to set up a USB device (enumeration)
- ☐ Initiated by the host
- ☐ Packet length
  - 8 bytes – low speed devices
  - 64 bytes – full speed devices
  - 8, 16, 32, or 64 bytes - high speed devices
- ☐ A Control Transfer can have up to three stages
  - Setup Stage
  - Data Stage
  - Status Stage

# Universal Serial Bus - USB 2.0

---

## ◆ Control Transfers

### □ Setup Stage

#### ▪ Comprises 3 packets

##### o Setup token

- contains address and endpoint number

##### o Data packet

- always has PID of type Data0
- Contains a setup packet with the type of request

##### o Handshake

- Acknowledge successful receipt or indicate an error
- If the function successfully receives the setup data (CRC and PID etc are okay), it responds with ACK
- *Otherwise, it ignores the data and doesn't send a handshake*
- Functions cannot issue STALL or NAK in response to a setup packet

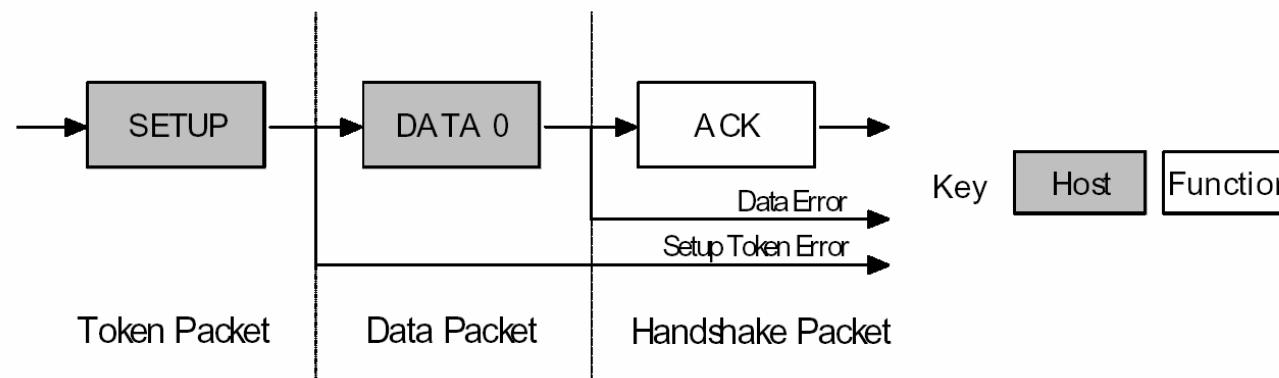


# Universal Serial Bus - USB 2.0

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## ◆ Control Transfers

### □ Setup Stage



# Universal Serial Bus - USB 2.0

---

## ◆ Control Transfers

### □ Data Stage (optional)

- The amount of data to be sent in this stage has been specified in the previous setup stage (in the setup packet)
- If the amount of data exceeds the maximum packet size, data is sent in multiple transfers (each of the maximum size except the last one)
- Two scenarios
  - **IN** direction
  - **OUT** direction

# Universal Serial Bus - USB 2.0

---

## ◆ Control Transfers

### □ Data Stage (optional)

- **IN** token issued when the host is ready to receive control data
  - If the function receives the **IN** token with an error, it ignores the packet
  - If the function receives the token correctly, it responds with either
    - **DATA** packet (with the control data to be sent), or
    - **STALL** packet (endpoint has had an error), or
    - **NAK** packet (endpoint if working but has no data to send yet)

# Universal Serial Bus - USB 2.0

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## ◆ Control Transfers

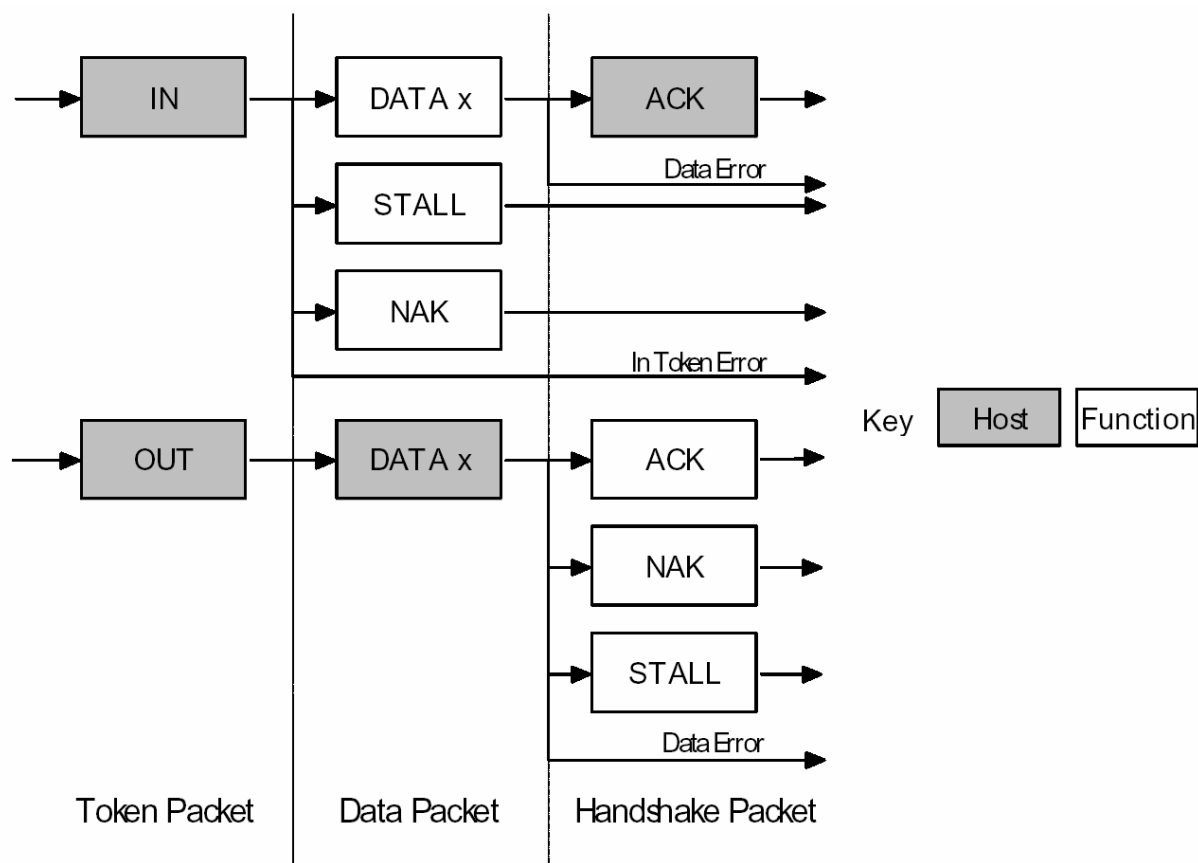
### □ Data Stage (optional)

- **OUT** token issued when the host needs to send the device a control data packet
- Followed by a **DATA** packet with the control data as a payload
  - If the function receives either the **OUT** token or the **DATA** packet with an error, it ignores the packet
  - Otherwise, it issues one of three responses
    - **ACK** (if the function receives the data correctly - successful transfer to the endpoint buffer)
    - **NAK** (if the buffer was not empty e.g. still processing previous packet)
    - **STALL** (endpoint has had an error)

# Universal Serial Bus - USB 2.0

## ◆ Control Transfers

- Data Stage (optional)



# Universal Serial Bus - USB 2.0

---

## ◆ Control Transfers

### □ Status Stage

- Reports the status of the overall request
- Varies depending on the direction of transfer
- Status reporting is always performed by the function
- **IN**
  - o If the host sent **IN** token during the data stage to receive data, the host must acknowledge successful receipt of the data
  - o Done by host sending an **OUT** token followed by a zero length data packet
  - o Function can now report status
    - **ACK** (success; ready for another command)
    - **STALL** (error)
    - **NAK** (still processing)

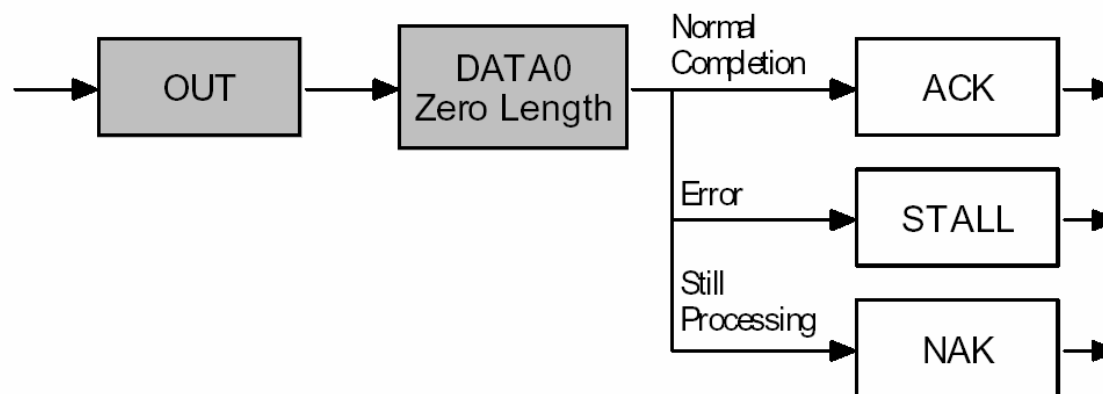
# Universal Serial Bus - USB 2.0

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## ◆ Control Transfers

### □ Status Stage

#### ▪ IN



# Universal Serial Bus - USB 2.0

---

## ◆ Control Transfers

### □ Status Stage

#### ▪ **OUT**

- o If the host sent **OUT** token(s) during the data stage to transmit data, the function must acknowledge successful receipt of the data
- o Done by host sending an **IN** token
- o Function responds:
  - **DATA0** (success; ready for another command)
  - **STALL** (error)
  - **NAK** (still processing; host has to retry status stage later)
- o Host completes the stage by sending an **ACK**

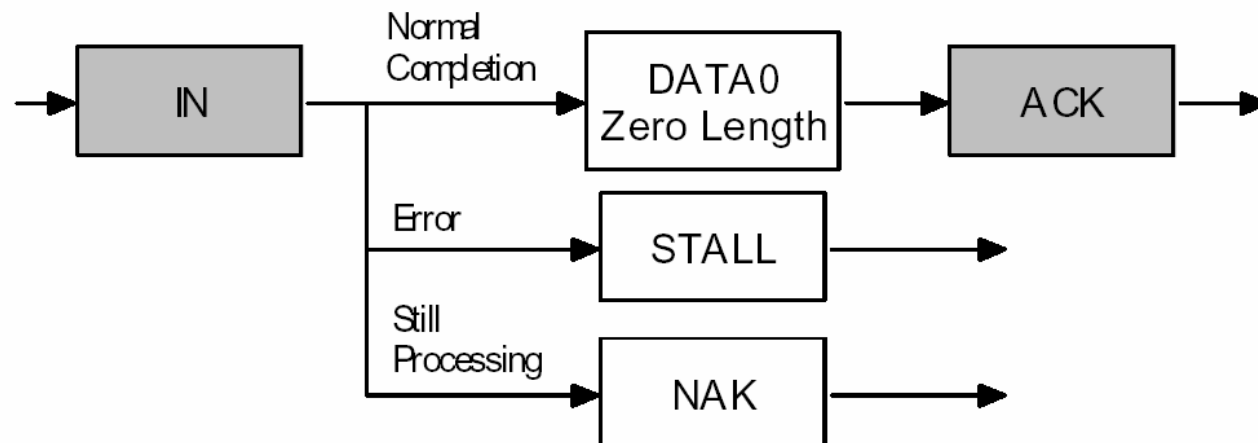


# Universal Serial Bus - USB 2.0

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## ◆ Control Transfers

- Status Stage
  - **OUT**



# Universal Serial Bus - USB 2.0

---

## ◆ Control Transfer Example

- ❑ Host wants to request a device descriptor during enumeration
- ❑ Host sends **Setup** token ( $\Rightarrow$  next packet is a setup packet)
  - Address field will hold the address of the device / function
  - Endpoint should be zero ( $\Rightarrow$  default pipe)
- ❑ Host sends **DATA0** packet
  - 8 byte payload (containing device descriptor request; ch. 9 USB Specification)
- ❑ Device send **ACK** (or ignores it, if there is an error)
- ❑ This is the first USB transaction (setup stage): the device now decodes the 8 bytes received, determines it is a device descriptor request, & then attempt to send the descriptor. This is the next transaction (next slide)

# Universal Serial Bus - USB 2.0

---

1. Setup Token

Sync	PID	ADDR	ENDP	CRC5	EOP
------	-----	------	------	------	-----

Address & Endpoint Number

2. Data0 Packet

Sync	PID	Data0	CRC16	EOP
------	-----	-------	-------	-----

Device Descriptor Request

3. Ack Handshake

Sync	PID	EOP
------	-----	-----

Device Ack. Setup Packet

# Universal Serial Bus - USB 2.0

---

## ◆ Control Transfer Example

- ❑ Assume the maximum payload size is 8 bytes
- ❑ Host now send **IN** token (tell device it can now send data)
- ❑ Device descriptor is sent in chunks of 8 bytes:
  - Device sends **DATA1** packet
  - Host sends **ACK** handshake
  - Host sends **IN** token
  - Device sends **DATA0** packet
  - Host sends **ACK** handshake
  - Host sends **IN** token
  - Device sends **DATA1** packet (last 8 bytes of descriptor)
  - Host send **ACK** handshake
- ❑ This is the second USB transaction (data stage)

# Universal Serial Bus - USB 2.0

---

1. In Token	Sync	PID	ADDR	ENDP	CRC5	EOP	Address & Endpoint Number
2. Data1 Packet	Sync	PID	Data1		CRC16	EOP	First 8 bytes of Device Descriptor
3. Ack Handshake	Sync	PID	EOP				Host Acknowledges Packet

1. In Token	Sync	PID	ADDR	ENDP	CRC5	EOP	Address & Endpoint Number
2. Data0 Packet	Sync	PID	Data0		CRC16	EOP	Second 8 bytes of Device Desc
3. Ack Handshake	Sync	PID	EOP				Host Acknowledges Packet

1. In Token	Sync	PID	ADDR	ENDP	CRC5	EOP	Address & Endpoint Number
2. Data1/0 Packet	Sync	PID	Data0/1		CRC16	EOP	Last 8 bytes of Device Descriptor
3. Ack Handshake	Sync	PID	EOP				Host Acknowledges Packet

# Universal Serial Bus - USB 2.0

---

## ◆ Control Transfer Example

- ☐ Host sends **OUT** Token
- ☐ Host sends a zero length **DATA1** packet (overall transaction is successful)
- ☐ Device sends **ACK** handshake
- ☐ This is the third USB transaction (status stage)

# Universal Serial Bus - USB 2.0

---

1. Out Token

Sync	PID	ADDR	ENDP	CRC5	EOP
------	-----	------	------	------	-----

Address & Endpoint Number

2. Data1 Packet

Sync	PID	Data1	CRC16	EOP
------	-----	-------	-------	-----

Zero Length Packet

3. Ack Handshake

Sync	PID	EOP
------	-----	-----

Function Ack. Entire Transactions

# Universal Serial Bus - USB 2.0

---

## ◆ Interrupt Transfers

- ❑ Under USB, if a device requires the attention of the host, it must wait until the host polls it before it can report it needs attention!
- ❑ Characteristics of Interrupt Transfers
  - Guaranteed latency
  - Stream pipe – unidirectional
  - Error detection and next period retry
- ❑ Typically non-periodic small device-initiated communication requiring bounded latency
- ❑ Maximum data payload
  - 8 bytes – low speed devices
  - 64 bytes – full speed devices
  - 1024 bytes - high speed devices



# Universal Serial Bus - USB 2.0

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## ◆ Interrupt Transfers

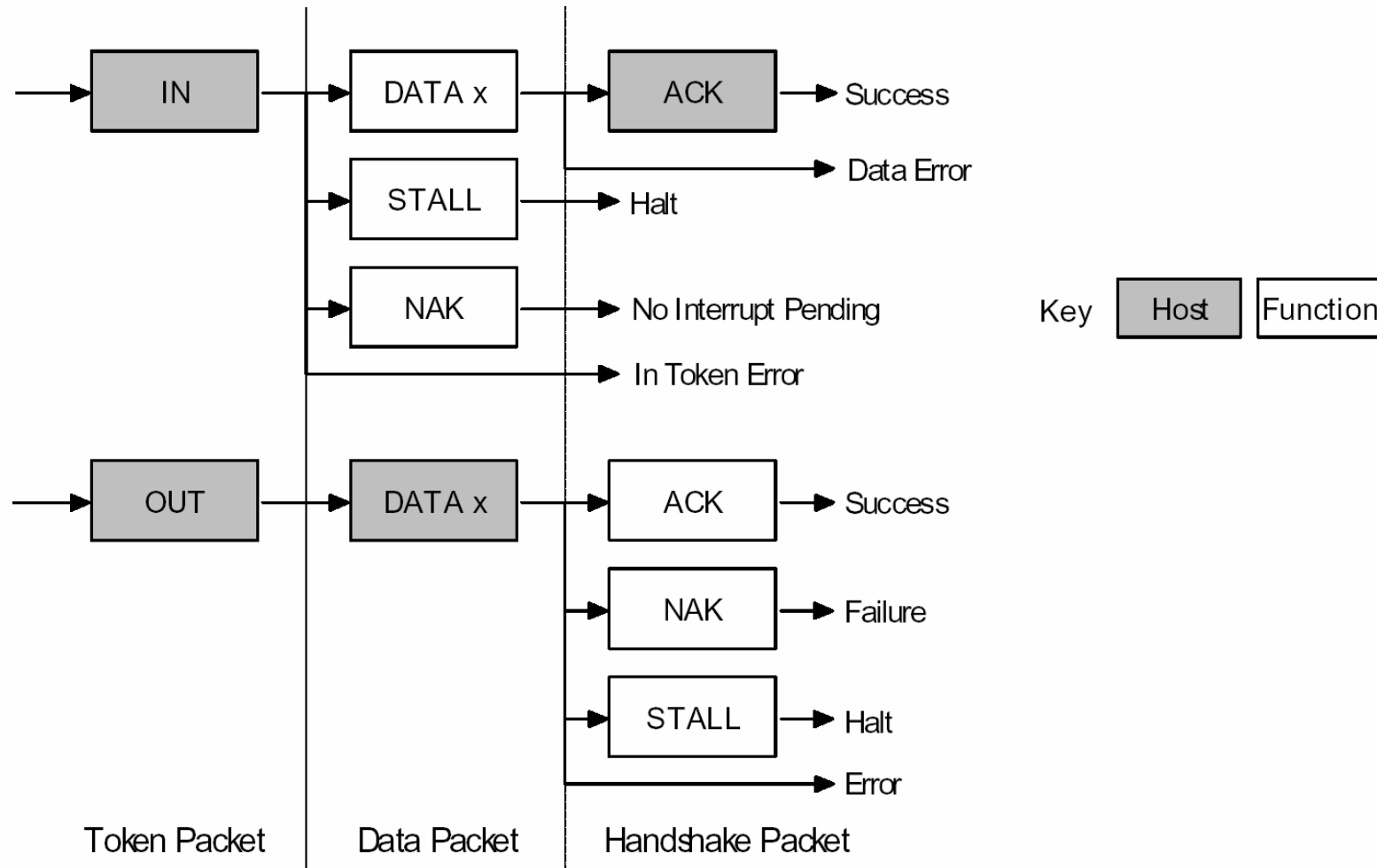
### ☐ Interrupt IN Transaction

- Host will periodically poll the interrupt endpoint
- Polling rate is specified in the endpoint descriptor
- Host sends an IN token, the function respond with a DATA packet, and the host acknowledges

### ☐ Interrupt OUT Transaction

- When the host wants to send the device interrupt data
- Sends OUT token, followed by a data packet, and device acknowledges

# Universal Serial Bus - USB 2.0



# Universal Serial Bus - USB 2.0

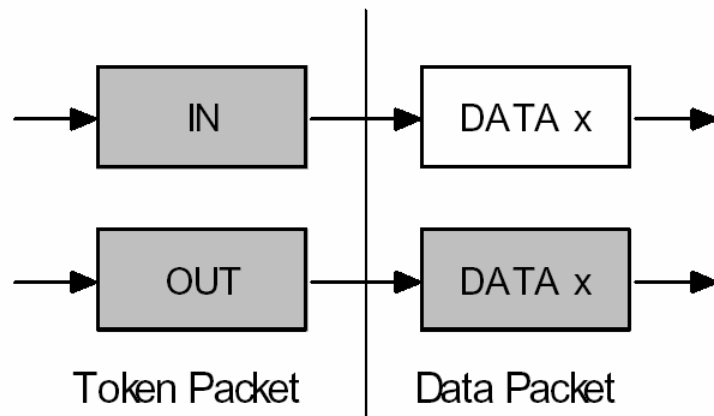
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## ◆ Isochronous Transfers

- ☐ Occur continuously and periodically
- ☐ Typically contain time sensitive information
  - Audio
  - Video
- ☐ Characteristics of Isochronous Transfers
  - Guaranteed access to USB bandwidth
  - Bounded latency
  - Stream pipe – unidirectional
  - Error detection via CRC, but no retry or guarantee of delivery
  - Full and high-speed modes only
- ☐ Maximum data payload is specified in the endpoint descriptor
  - 1023 bytes – full speed devices
  - 1024 bytes - high speed devices
  - Amount of data can be less than the pre-negotiated size and may vary from transaction to transaction
- ☐ No handshaking
- ☐ No error reports or STALL/HALT conditions

# Universal Serial Bus - USB 2.0

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# Universal Serial Bus - USB 2.0

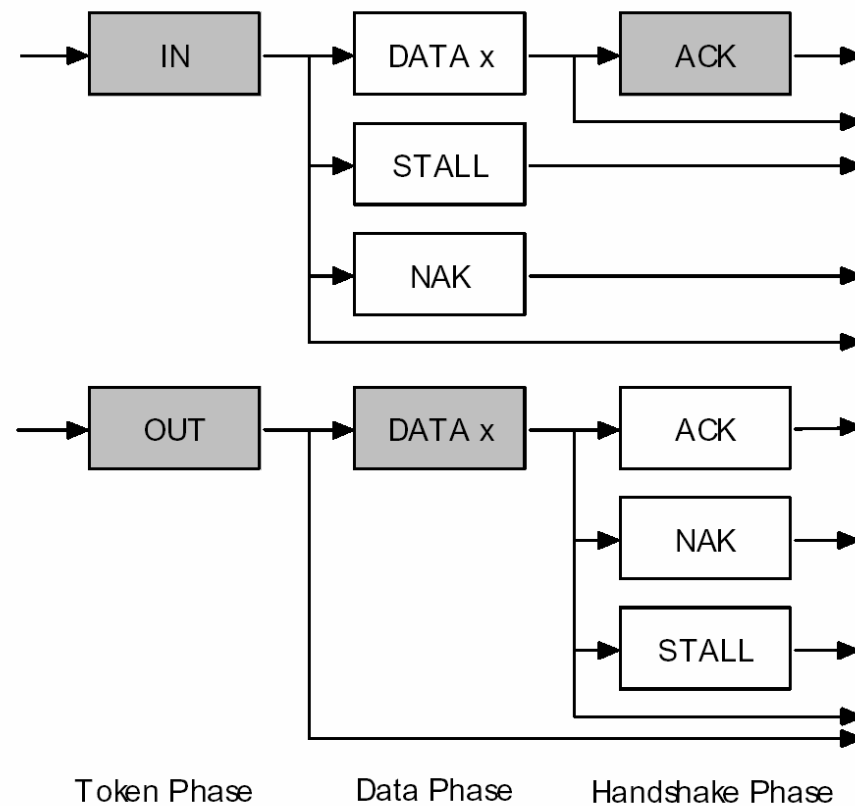
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## ◆ Bulk Transfers

- ❑ Used for large bursty data transfer (e.g. printer, scanner)
- ❑ Characteristics of Bulk Transfers
  - Error detection via CRC, with guarantee of delivery
  - No guarantee of bandwidth or minimum latency
  - Stream pipe – unidirectional
  - Full and high-speed modes only
- ❑ Maximum data payload
  - 8, 16, 32, 64 bytes – full speed devices
  - 512 bytes - high speed devices
  - Amount of data can be less than the maximum packet size (zero padding not needed)

# Universal Serial Bus - USB 2.0

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# Universal Serial Bus - USB 2.0

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## ◆ Bandwidth Management

- ☐ The host is responsible for managing bandwidth
- ☐ Achieved at enumeration when configuring isochronous and interrupt endpoints (and throughout the operation of the bus)
- ☐ USB Specification Limits:
  - Full speed bus:  $\leq 90\%$  allocated to periodic transfers
  - High speed bus:  $\leq 80\%$  allocated to periodic transfers
  - Control transfer get allocated from the remainder
  - Bulk transfers get what's left!

# Universal Serial Bus - USB 2.0

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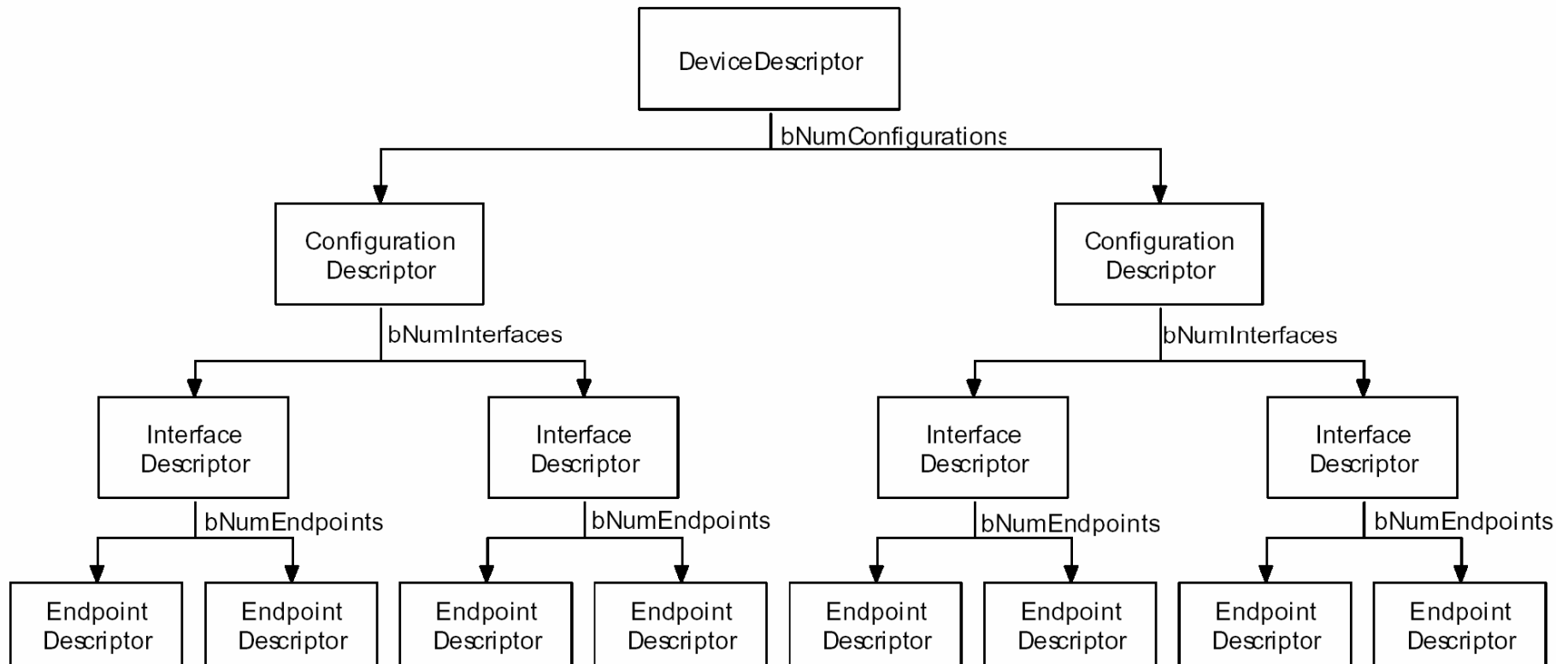
## ◆ USB Descriptors

- ❑ All USB devices define a hierarchy of descriptors
  - What the device is
  - What version of USB it supports
  - How many ways it can be configured
  - Number of end-points and their types
  
- ❑ Common descriptors
  - Device descriptors
  - Configuration descriptors
  - Interface Descriptors
  - Endpoint descriptors
  - String descriptors



# Universal Serial Bus - USB 2.0

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# Universal Serial Bus - USB 2.0

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## ◆ USB Descriptors

### ☐ Device Descriptor

- Only one of these
- USB revision compliance
- Product id and vendor id (needed to load the drivers)
- Number of configurations supported

### ☐ Configuration Descriptor

- Amount of power used
- Self- or bus-powered
- Number of interfaces
- Very few devices have more than one configuration

# Universal Serial Bus - USB 2.0

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## ◆ USB Descriptors

### ☐ Interface Descriptor

- Groups endpoints into functional groups
- Each functional group corresponds to a feature of the device
- For example, a fax/scanner/printer device might have three groups (and, hence, three interface descriptors)

### ☐ Endpoint Descriptors

- Describe endpoints other than endpoint zero
- Endpoint zero is always assumed to be a control endpoint

### ☐ String Descriptors

- Provide human readable information
- Optional

# Universal Serial Bus - USB 2.0

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## ◆ The Setup Packet

- ☐ Every USB device must respond to setup packets on the default pipe
- ☐ Used for detection and configuration of the device
  - Setting USB device's address
  - Requesting device descriptor
  - Checking the status of an endpoint
- ☐ All requests have to be processed within set time limits
- ☐ Setup Packet is 8 bytes long

# Universal Serial Bus - USB 2.0

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Offset	Field	Size	Value	Description
0	bmRequestType	1	Bit-Map	<b>D7 Data Phase Transfer Direction</b> 0 = Host to Device 1 = Device to Host <b>D6..5 Type</b> 0 = Standard 1 = Class 2 = Vendor 3 = Reserved <b>D4..0 Recipient</b> 0 = Device 1 = Interface 2 = Endpoint 3 = Other 4..31 = Reserved
1	bRequest	1	Value	Request
2	wValue	2	Value	Value
4	wIndex	2	Index or Offset	Index
6	wLength	2	Count	Number of bytes to transfer if there is a data phase

# Universal Serial Bus - USB 2.0

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## ◆ The Setup Packet

- bRequest field determines the request being made
  - Standard device requests (section 9.4 of the USB specification)
    - Must be implemented for every USB device
      - GET\_STATUS
      - SET\_ADDRESS
      - 6 more
  - Class requests (e.g. communications class, mass storage class)
  - Vendor defined requests

# Universal Serial Bus - USB 2.0

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## ◆ Enumeration

- ☐ The process of determining
  - what device has just been connected to the bus
  - The power requirements
  - Number and type of endpoints
  - Class of product
- ☐ Host then assigns an address to the device
- ☐ Host also enables a configuration allowing the device to transfer data on the bus
- ☐ See section 9.1.2 of the USB specification

# Universal Serial Bus - USB 2.0

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- ◆ We have concentrated on the USB interface from the point of view of the device or function
- ◆ However, something has to manage all the transactions on the bus ... this is the job of the host
- ◆ We mentioned three USB host controller specifications
  - ❑ UHCI (Universal Host Controller Interface)
    - Developed by Intel
    - More work is done in software
    - Hence cheaper hardware
  - ❑ OHCE (Open Host Controller Interface)
    - Developed by Compaq, Microsoft, and National Semiconductor
    - More work done in hardware
    - Hence simpler software!
  - ❑ USB 2.0
    - EHCI (Enhanced Host Controller Interface)



# Universal Serial Bus - USB 2.0

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- ◆ The job of each of these interfaces is to schedule transactions on the bus, based on the information provided by each function/device at enumeration
  - ❑ Meet the needs of all isochronous, interrupt, control, and bulk transfers
- ◆ Transactions are generated by cycling through a list of frames (one frame every 1ms), each frame comprising a list of transfer descriptors
  - ❑ Transfer descriptor
    - USB device address
    - Type of transaction to be performed
    - Transfer size
    - Speed of transaction
    - Location of memory data buffer (where data will be read from or written to)

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# The IEEE-488 Bus

# IEEE-488 Bus

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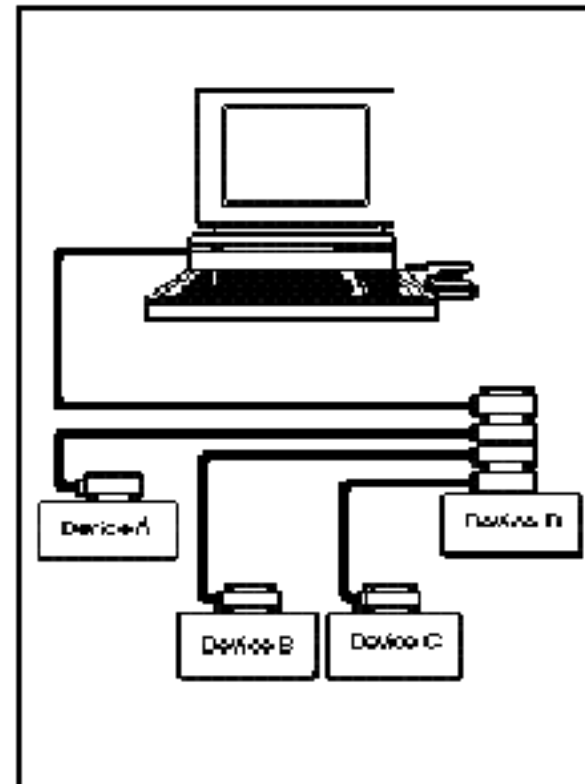
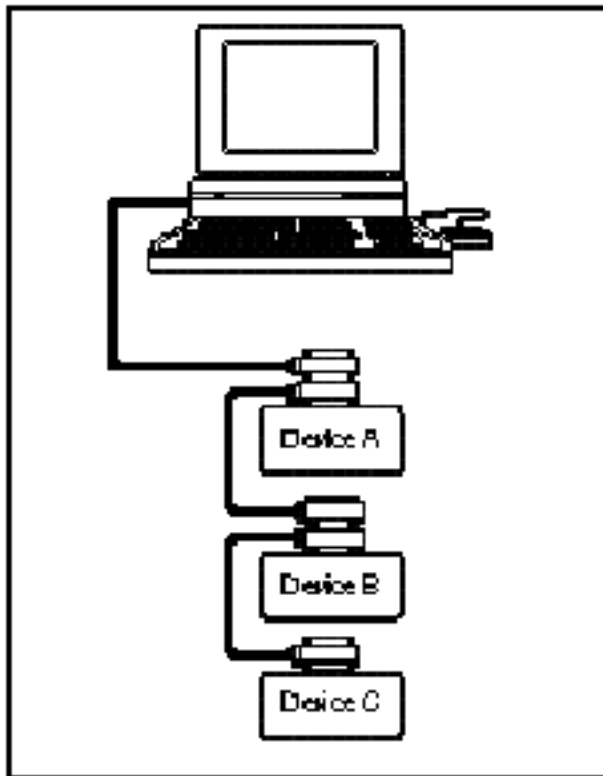
- ◆ Digital Interface for Programmable Instrumentation
- ◆ Also known as the General Purpose Instrument Bus (GPIB)
  - ❑ Provides a way of interconnecting a microcomputer controller with a large range of test and measuring instruments
    - Signal generators, oscilloscopes, network analysers ...
  - ❑ The objective of the IEEE-488 bus is to enable different instruments using different data transfer rates, different message lengths, and different capabilities to communicate with each other on a single bus
  - ❑ Allows a PC to supervise operation of instruments, and collect & process data they provide

(Based in part on information abstracted from: [http://www.interfacebus.com/Design\\_Connector\\_GPIB.html](http://www.interfacebus.com/Design_Connector_GPIB.html) and <http://www.hit.bme.hu/people/papay/edu/GPIB/tutor.htm> )

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# IEEE-488 Bus

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Supports both star and daisy-chain topologies

# IEEE-488 Bus

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- ◆ In 1965, Hewlett-Packard designed the Hewlett-Packard Interface Bus (**HP-IB**) to connect their line of programmable instruments to their computers
- ◆ Because of its high transfer rates (nominally 1 Mbytes/s), this interface bus quickly gained popularity
- ◆ It was later accepted as IEEE Standard 488-1975, and has evolved to ANSI/IEEE Standard **488.1**-1987
- ◆ Today, the name **General Purpose Interface Bus (GPIB)** is more widely used than HP-IB.
- ◆ ANSI/IEEE **488.2**-1987 strengthened the original standard by defining precisely how controllers and instruments communicate

# IEEE-488 Bus

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- ◆ The IEEE-488 interface bus is an 8 bit wide byte serial, bit parallel interface system which incorporates:
  - ❑ 5 control lines (interface management)
  - ❑ 3 handshake lines
  - ❑ 8 bi-directional data lines.
- ◆ The entire bus consists of 24 lines, with the remaining lines occupied by ground wires.
- ◆ Uses negative logic with standard TTL levels
  - ❑ When DAV is true, for example, it is a TTL low level ( $\leq 0.8$  V)
  - ❑ When DAV is false, it is a TTL high level ( $\geq 2.0$  V)
- ◆ The maximum data transfer rate is determined by a number of factors, but is assumed to be 1Mb/s.

# IEEE-488 Bus

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- ◆ Devices exist on the bus in any one of 3 general forms:
  1. Controller
  2. Talker
  3. Listener
  
- ◆ A single device may incorporate all three options, although only one option may be active at a time.

# IEEE-488 Bus

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## ◆ CONFIGURATION REQUIREMENTS

- ☐ To achieve the high data transfer rate for which the GPIB was designed, the physical distance between devices and the number of devices on the bus are limited
- ☐ A maximum separation of 4 m between any two devices and an average separation of **2 m** over the entire bus
- ☐ A maximum total cable length of **20 m**
- ☐ No more than **15 device** loads connected to each bus, with no less than two-thirds powered on



# IEEE-488 Bus

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- ◆ The Controller makes the determination as to which device becomes active on the bus
- ◆ The GPIB can handle only one 'active' controller on the bus, although it may pass operation to another controller.
- ◆ Any number of active listeners can exist on the bus with an active talker as long as no more than 15 devices are connected to the bus
- ◆ The controller determines which devices become active by sending interface messages over the bus to a particular instrument

# IEEE-488 Bus

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- ◆ Each individual device is associated with a 5 bit BCD code which is unique to that device.
  - ❑ By using this code, the controller can coordinate the activities on the bus and the individual devices can be made to talk, listen (un-talk, un-listen) as determined by the controller.
  - ❑ A controller can only select a particular function of a device, if that function is incorporated within the device; for example a 'listen' only device can not be made to talk to the controller.
- ◆ The Talker sends data to other devices
- ◆ The Listener receives the information from the Talker

# IEEE-488 Bus

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- ◆ Device dependent messages are moved over the GPIB in conjunction with the data byte transfer control lines
- ◆ These three lines (DAV, NRFD, and NDAC) are used to form a three wire 'interlocking' handshake which controls the passage of data
  - ❑ The active talker would control the 'DAV' line (Data Valid)
  - ❑ The listener(s) would control the 'NRFD' (Not Ready For Data), and the 'NDAC' (Not Data Accepted) line

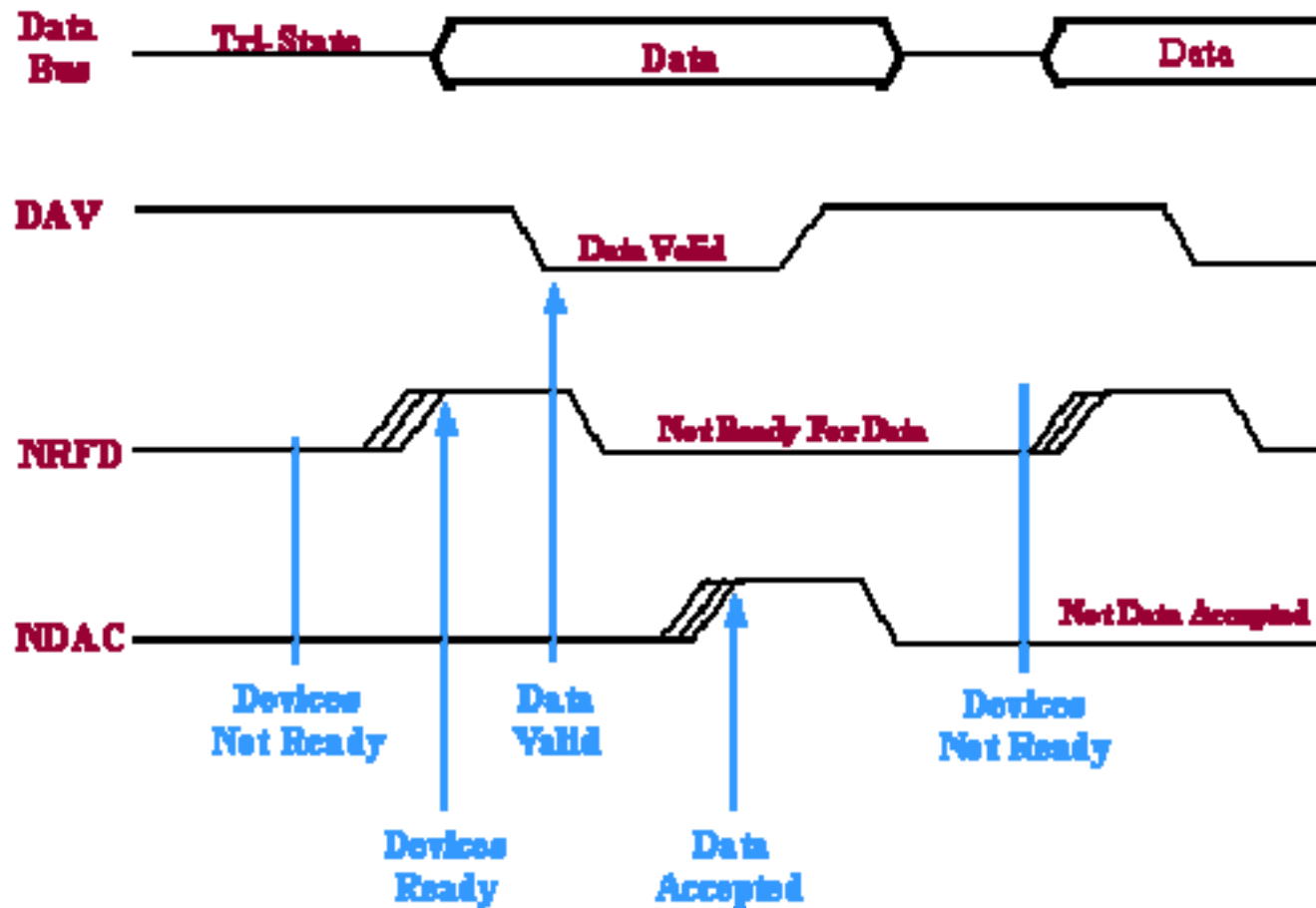
# IEEE-488 Bus

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- ◆ In the steady state mode the talker will hold 'DAV' high (no data available) while the listener would hold 'NRFD' high (ready for data) and 'NDAC' low (no data accepted)
- ◆ After the talker placed data on the bus it would then take 'DAV' low (data valid)
- ◆ The listener(s) would then send 'NRFD' low and send 'NDAC' high (data accepted)
- ◆ Before the talker lifts the data off the bus, 'DAV' will be taken high signifying that data is no longer valid.

# IEEE-488 Bus

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# IEEE-488 Bus

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- ◆ If the 'ATN' line (attention) is high while this process occurs the information is considered data ( a device dependent message), but with the "ATN' line low the information is regarded as an interface message; such as listen, talk, un-listen or un-talk.
- ◆ The other five lines on the bus ('ATN' included) are the bus management lines
- ◆ These lines enable the controller and other devices on the bus to enable, interrupt, flag, and halt the operation of the bus

# IEEE-488 Bus

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- ◆ Most devices operate either via front panel control or GPIB control (REMOTE)
- ◆ While using the front Panel the device is in the Local state, when receiving commands via the GPIB, the device is in the Remote state
- ◆ The device is placed in the Remote state when ever the System Controller is reset or powered on; also, when the system controller sends out an Abort message
- ◆ In addition, if the device is addressed, it then enters the Remote state.

# IEEE-488 Bus

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Pin	Signal Name	Function	Pin	Signal Name	Function
1	DIO1	Data input/output bit 1	13	DIO5	Data input/output bit 5
2	DIO2	Data input/output bit 2	14	DIO6	Data input/output bit 6
3	DIO3	Data input/output bit 3	15	DIO7	Data input/output bit 7
4	DIO4	Data input/output bit 4	16	DIO8	Data input/output bit 8
5	EOI	End-of-identify	17	REN	Remote enable
6	DAV	Data valid	18	SHIELD	Ground (DAV)
7	NRFD	Not ready for data	19	SHIELD	Ground (NRFD)
8	NDAC	Not data accepted	20	SHIELD	Ground (NDAC)
9	IFC	Interface clear	21	SHIELD	Ground (IFC)
10	SRQ	Service request	22	SHIELD	Ground (SRQ)
11	ATN	Attention	23	SHIELD	Ground (ATN)
12	SHIELD	Chassis ground	24	SIGNAL GND	Signal ground



# IEEE-488 Bus

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- ◆ NRFD: Not Ready For Data
  - ☐ Part of a three wire handshake
  - ☐ Used to indicate a device is ready for data, active low
- ◆ DAV: Data Valid
  - ☐ Part of a three wire handshake
  - ☐ Used to indicate valid data on the bus, active low
- ◆ NDAC: Not Data Accepted
  - ☐ Part of a three wire handshake
  - ☐ Used to indicate a device has not yet accepted data, active low

# IEEE-488 Bus

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## ◆ ATN: Attention

- ☐ When low (true) the system places all devices in Command Mode
- ☐ When high (false) the system places all devices in the Data Mode
- ☐ In Command Mode the Controller passes data to devices,
- ☐ In Data Mode the Talker passes data to the Listener.
- ☐ All devices must monitor the ATN line and respond within 200nS.

# IEEE-488 Bus

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- ◆ EIO: End or Identify
  - ❑ Indicates the last data transfer of a multi-byte sequence or used by the system controller to indicate a Parallel Poll to a device (in conjunction with the ATN line)

# IEEE-488 Bus

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## ◆ IFC: Interface Clear;

- ☐ Used only by the system controller to halt current operations
- ☐ Placing all devices in an idle state
- ☐ All talkers are set to Un-talk and all Listeners are set to Un-listen
- ☐ Serial Poll is disabled
- ☐ All devices must monitor IFC, and respond within 100uS

# IEEE-488 Bus

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## ◆ REN: Remote Enable

- ☐ Used by the system controller to place devices in programming mode
- ☐ All Remote capable Listeners are set to remote operation ( if they were addressed to Listen), when REN is true (low)
- ☐ When false, devices are set to Local control
- ☐ All device must monitor the REN line, and respond within 100uS.

# IEEE-488 Bus

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- ◆ SRQ: Service Request
  - ☐ Used by any device to indicate that a device needs service
  - ☐ Any device may use this line
  - ☐ Normally used as an interrupt line
  - ☐ The controller may mask this line, in favor of Polling
  - ☐ The SRQ line is cleared by a Serial Poll

# IEEE-488 Bus

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- ◆ DIO1 to DIO8: Data Input-Output bus
  - ❑ Bi-directional, Bite Serial-Byte Parallel.
  - ❑ Per data line, bits are serial, ASCII data is sent as parallel data.  
Standard data is 7 bit ASCII
  - ❑ But no coding format is defined with IEEE-488.

# IEEE-488 Bus

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- ◆ **IEEE 488.2 AND SCPI**
- ◆ The SCPI and IEEE 488.2 standards addressed the limitations and ambiguities of the original IEEE 488 standard
- ◆ **IEEE 488.2** makes it possible to design more compatible and productive test systems



# IEEE-488 Bus

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- ◆ The ANSI/IEEE Standard 488-1975, now called **IEEE 488.1**, greatly simplified the interconnection of programmable instrumentation by clearly defining mechanical, electrical, and hardware protocol specifications
- ◆ For the first time, instruments from different manufacturers were interconnected by a standard cable
- ◆ Although this standard went a long way towards improving the productivity of test engineers, the standard did have a number of shortcomings
- ◆ Specifically, IEEE 488.1 did not address data formats, status reporting, message exchange protocol, common configuration commands, or device-specific commands. As a result, each manufacturer implemented these items differently, leaving the test system developer with a formidable task.

# IEEE-488 Bus

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- ◆ **IEEE 488.2** enhanced and strengthened IEEE 488.1 by standardizing data formats, status reporting, error handling, Controller functionality, and common commands to which all instruments must respond in a defined manner
- ◆ By standardizing these issues, IEEE 488.2 systems are much more compatible and reliable
- ◆ The IEEE 488.2 standard focuses mainly on the software protocol issues and thus maintains compatibility with the hardware-oriented IEEE 488.1 standard

# IEEE-488 Bus

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National Instruments: <http://sine.ni.com/apps/we/nio.vp?cid=1230&lang=US>

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# IEEE-488 Bus

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# IEEE-488 Bus

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# Digital Input/Output

# Digital I/O

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- ◆ Digital input is the simplest form of inputting signals created by external devices to a computer
- ◆ Digital output is the simplest form of outputting from a computer to external control devices
- ◆ Digital input/output (I/O) provides for the inputting of digital logic signals to the computer from 'real-world' and outputting of compatible logic signals to the 'real-world'
- ◆ Typically uses an I/O Data Acquisition & Control (DA&C) board mapped on the PC bus to connect to switches, actuators, sensors, relays, lights, ...

# Digital I/O

---

- ◆ Digital I/O can be achieved using the Programmable Peripheral Interface (PPI) device
  - ❑ Intel 8255 IC
  - ❑ 40 pins
  - ❑ Normally mapped into the isolated I/O memory map of a PC
  - ❑ TTL voltage compatibility (5.5V maximum input)
  - ❑ Maximum current source / sink: 1mA



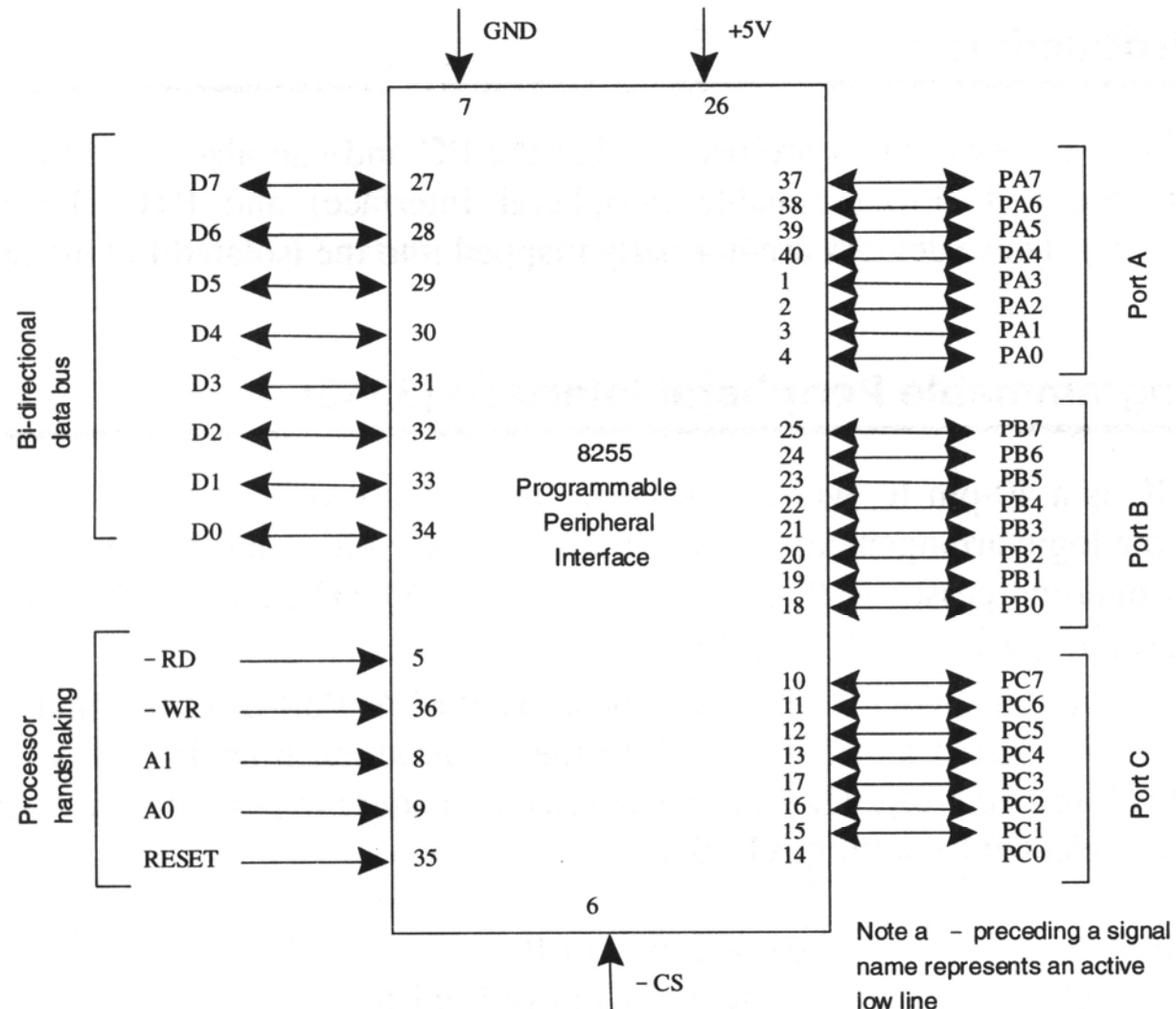
# Digital I/O

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- ◆ Connection to the system microprocessor is via the data bus (D0-D7), with handshaking lines:

- ☐  $\overline{RD}$  Read
- ☐  $\overline{WR}$  Write
- ☐ A1 Address line to select I/O port or control register
- ☐ A0
- ☐  $\overline{RESET}$  low input on RESET initializes the device
- ☐ CS chip select (low to activate the device)

# Digital I/O



# Digital I/O

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- ◆ 24 I/O lines, grouped into three groups of 8 bits
  - ❑ Port A (address 00 [A1A0] – BASE\_ADDRESS)
  - ❑ Port B (address 01 – BASE\_ADDRESS+1)
  - ❑ Port C (address 10 – BASE\_ADDRESS+2)
    - Port C is divided into two parts: upper and lower

# Digital I/O

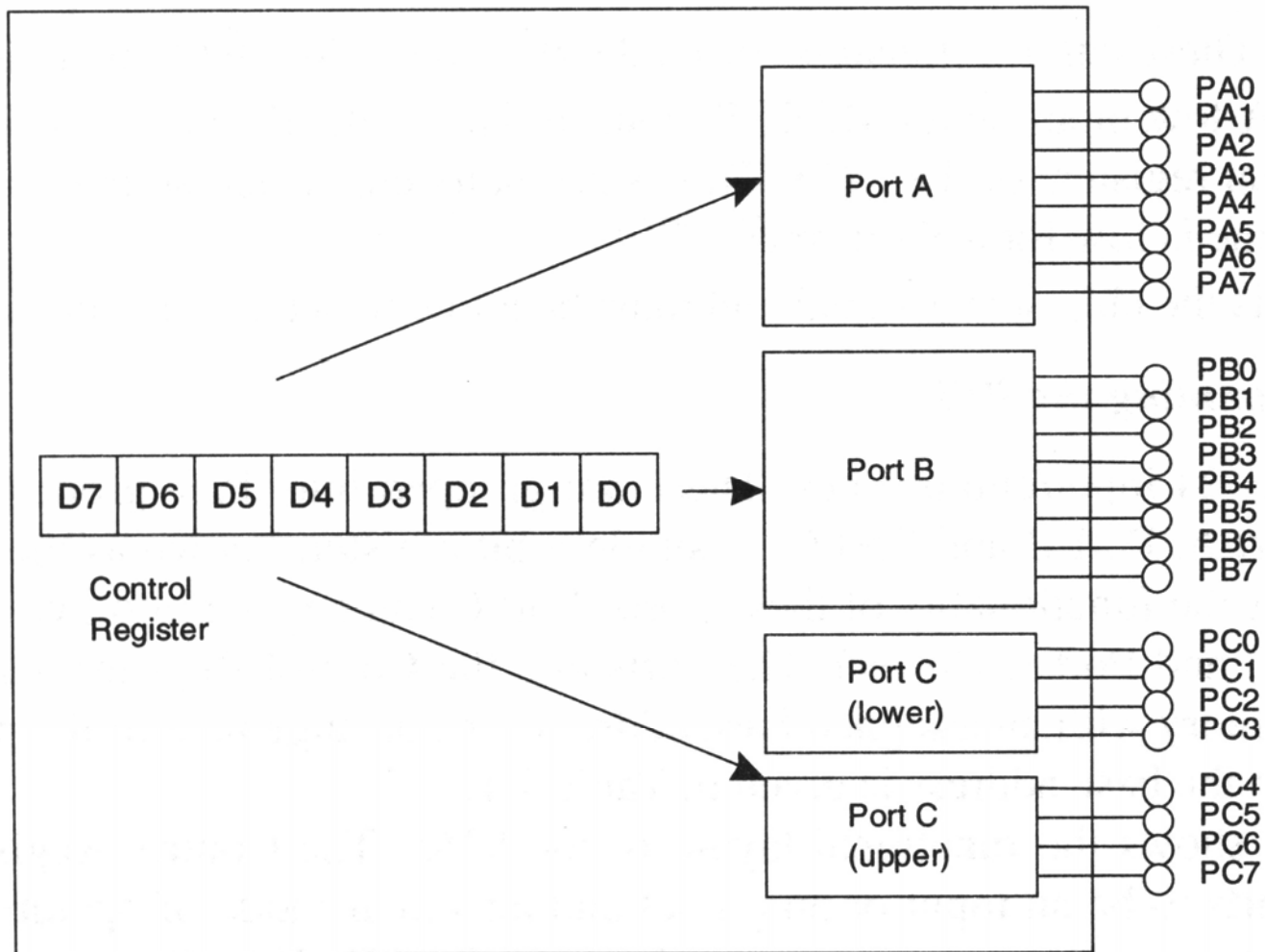
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◆ Functionality programmed using the 8-bit Control Register (address 11 – BASE\_ADDRESS+3)

<input type="checkbox"/>	D7	0	inactive
		1	active
<input type="checkbox"/>	D4	0	Port A is an output port
		1	Port A is an input port
<input type="checkbox"/>	D1	0	Port B is an output port
		1	Port B is an input port
<input type="checkbox"/>	D0	0	Port C lower is an output port
		1	Port C lower is an input port
<input type="checkbox"/>	D3	0	Port C upper is an output port
		1	Port C upper is an input port

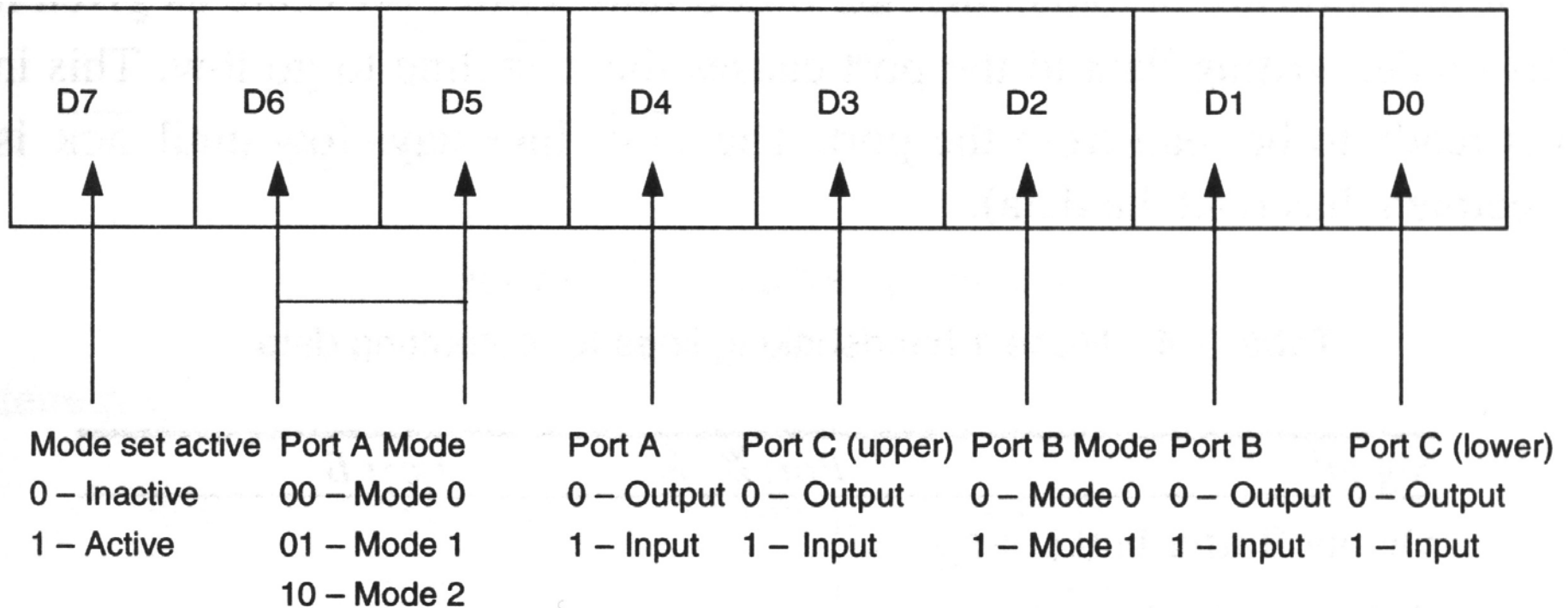
# Digital I/O

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# Digital I/O

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# Digital I/O

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- ◆ Port A can be operated in one of three modes (0, 1, 2) set by D6 & D5
  - Mode 0 (D6 D5 = 00)
    - Simplest mode: no handshaking
    - Bits on Port C can be programmed as inputs or outputs

# Digital I/O

---

- ◆ Port A can be operated in one of three modes (0, 1, 2) set by D6 & D5

- Mode 1 (D6 D5 = 01)

- Handshaking for synchronization (required when speed of devices differs significantly)
- Port C is used for the handshaking signals

- Input:

- o PC4       $\overline{\text{STB}}$       drive low to write data into Port A
- o PC5       $\overline{\text{IBF}}$       Input Buffer Full; automatically goes high when data has been written into Port A & remains high until data has been read

- Output:

- o PC7       $\overline{\text{OBF}}$       Output Buffer Full Port A, automatically goes low when data is written to the port
- o PC6       $\overline{\text{ACK}}$       drive low to acknowledge data has been read (and  $\overline{\text{OBF}}$  then goes high again)



# Digital I/O

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- ◆ Port A can be operated in one of three modes (0, 1, 2) set by D6 & D5

- Example: Port A input, mode 1 handshaking
  - External device placed data on Port A data lines
  - External device asserts  $\overline{STB}$  (i.e. 0, driven low)
  - Data is latched on port A register
  - 8255 asserts IBF (i.e. 1, driven high)
- Example: Port A output, mode 1 handshaking
  - CPU writes data to Port A
  - $\overline{OBF}$  is asserted (i.e. 0, driven low)
  - External device reads data on port A data lines
  - External device asserts  $\overline{ACK}$  (i.e. 0, driven low)
  - $\overline{OBF}$  is cleared (i.e. 1, driven high)

# Digital I/O

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- ◆ Port A can be operated in one of three modes (0, 1, 2) set by D6 & D5

- Mode 2 (D6 D5 = 10)

- Bi-directional I/O
- Port C is used for the handshaking signals
- Input:
  - o PC4       $\overline{\text{STB}}$       drive low to write data into Port A
  - o PC5       $\text{IBF}$       Input Buffer Full; automatically goes high when data has been written into Port A & remains high until data has been read
- Output:
  - o PC7       $\text{OBF}$       Output Buffer Full Port A, automatically goes low when data is written to the port
  - o PC6       $\text{ACK}$       drive low to acknowledge data has been read (and  $\text{OBF}$  then goes high again)

# Digital I/O

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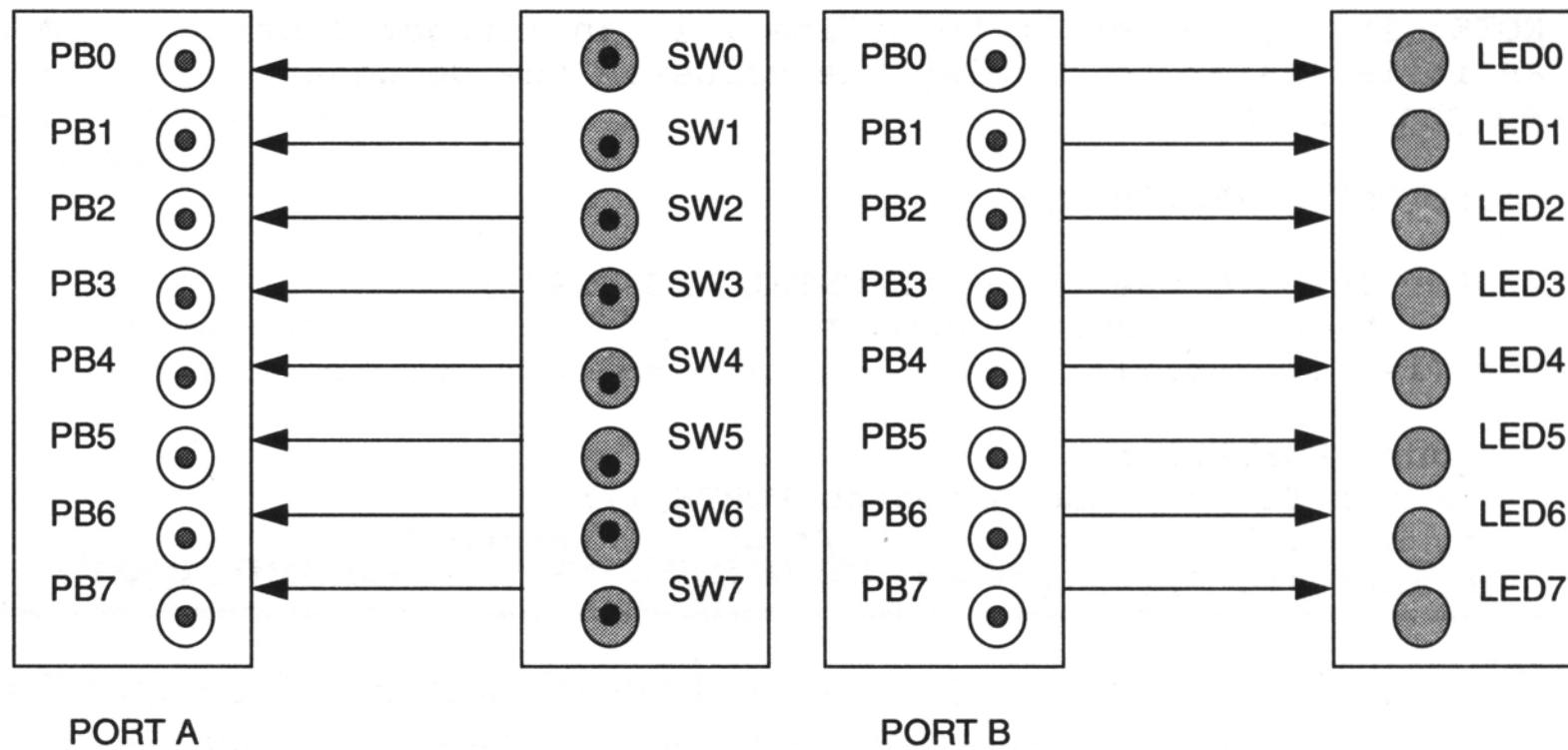
- ◆ Port B can be operated in one of two modes (0, 1) set by D2
  - ❑ Mode 0 (D2 = 0)
    - Simplest mode: no handshaking
    - Bits on Port C can be programmed as inputs or outputs
  - ❑ Mode 1 (D2 = 1)
    - Handshaking for synchronization (required when speed of devices differs significantly)
    - Port C is used for the handshaking signals
    - Input:

o PC2	$\overline{\text{STB}}$	drive low to write data into Port B
o PC1	IBF	Input Buffer Full; automatically goes high when data has been written into Port B & remains high until data has been read
    - Output:

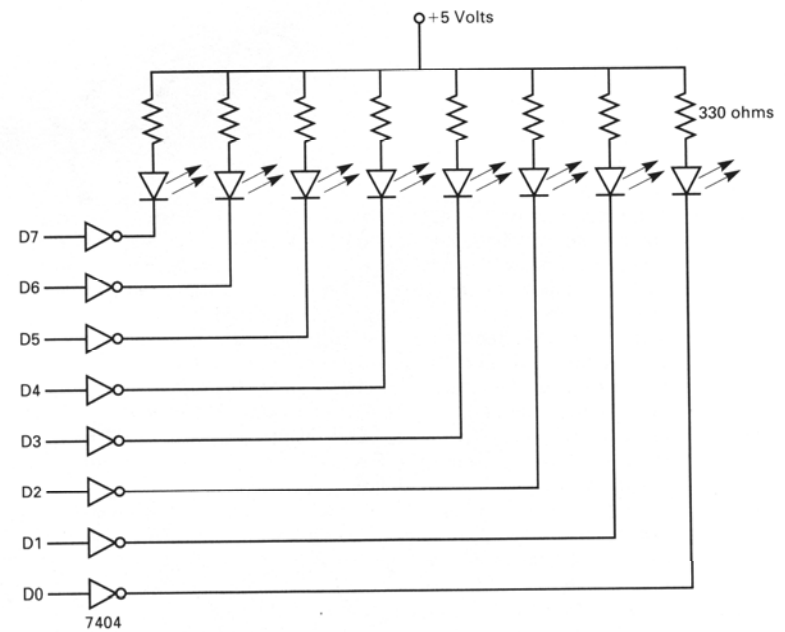
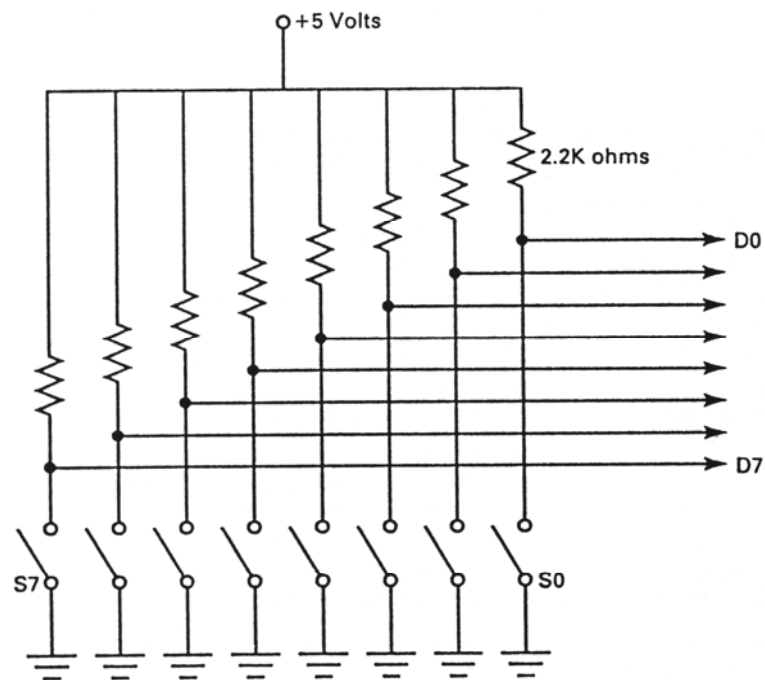
o PC1	$\overline{\text{OBF}}$	Output Buffer Full Port B, automatically goes low when data is written to the port
o PC2	$\overline{\text{ACK}}$	drive low to acknowledge data has been read (and $\overline{\text{OBF}}$ then goes high again)

# Digital I/O

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# Digital I/O



( Figures based on W. H. Rigby and T. Dalby, *Computer Interfacing* )  
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# Digital I/O

---

```
/*    ppi_2.c - read Port A and write same data to Port B                                */

#define BASE_ADDRESS    0x3B0  /* change this as required */
#define PORTA           BASE_ADDRESS
#define PORTB           (BASE_ADDRESS+1)
#define PORTC           (BASE_ADDRESS+2)
#define CNTRL_REG       (BASE_ADDRESS+3)
#include <stdio.h>
#include <time.h>        /* included for term()                                */
#include <dos.h>         /* included for inputb and outputb                      */

void          my_delay(int secs);

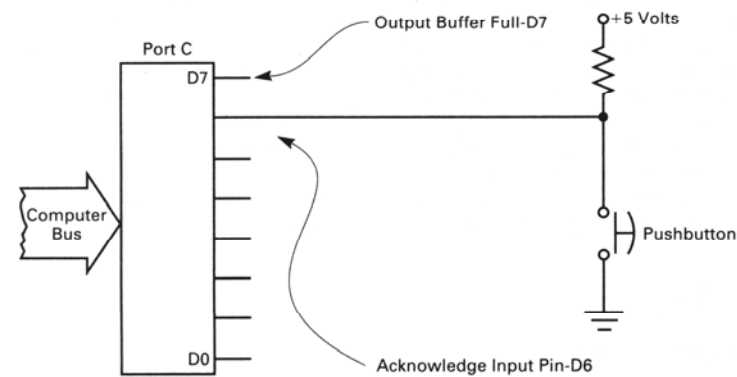
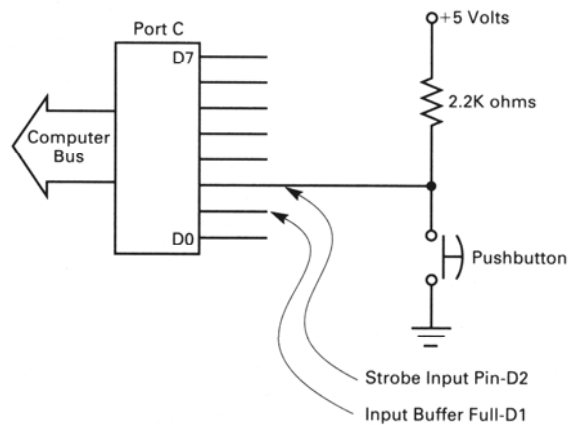
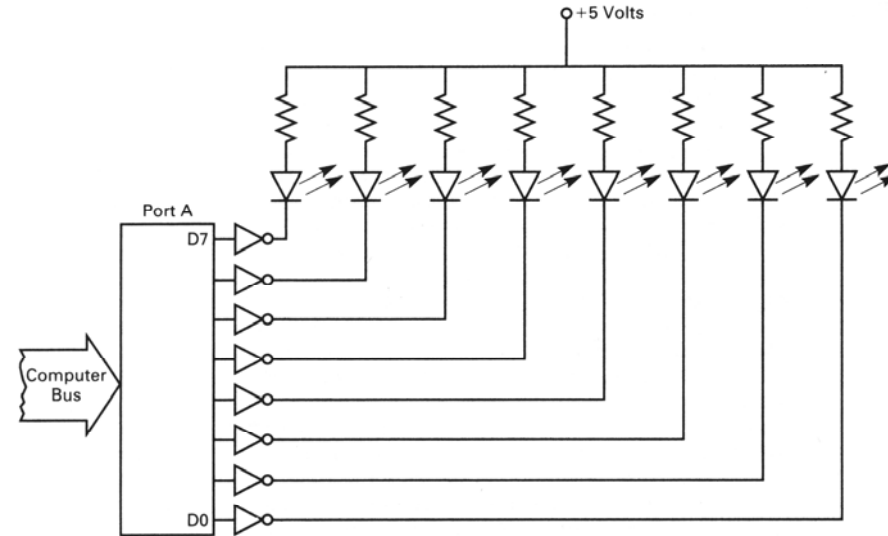
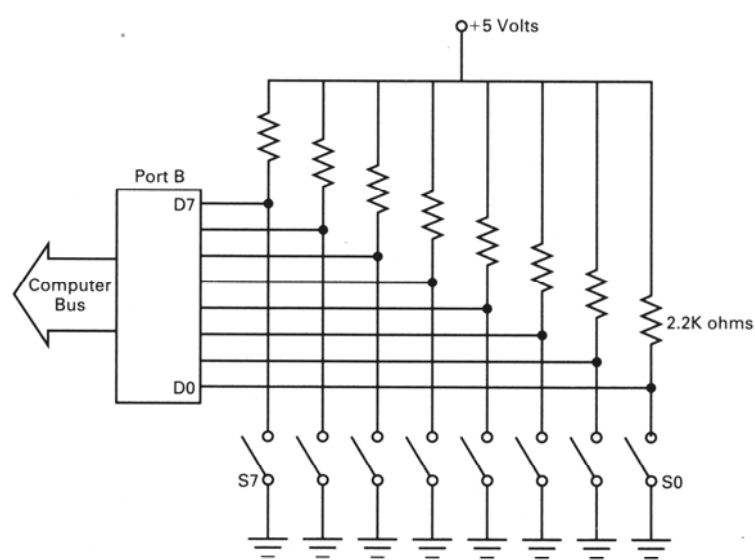
int main(void)
{
    unsigned char i=0;
    outportb(CNTRL_REG, 0x90); /* what does this do? */
    do
    {
        i = inportb(PORTA);
        outportb(PORTB, i);
        my_delay(1);          /* wait 1 second */
        printf("Input value is %d\n", i);
    } while (i != 0xff);
    return(0);
}
```

# Digital I/O

---

```
void my_delay(int secs)
{
    time_t oldtime, newtime;
    time(&oldtime);
    do {
        time(&newtime);
    } while ((newtime-oldtime)<secs);
}
```

# Digital I/O



( Figures based on W. H. Rigby and T. Dalby, *Computer Interfacing*)

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# Digital I/O

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## ◆ Optical Isolation

### □ Discrete Input Modules

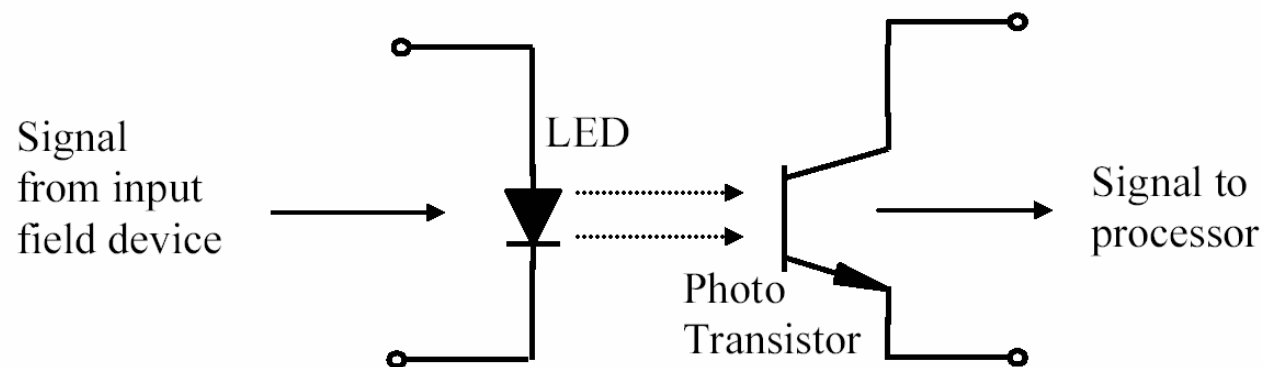
- Discrete I/O modules are installed as interfaces when discrete field devices are used to sense and control process variables. A discrete field device has only two functional states. These states may be on/off, open/closed, or true/false.
- Typical discrete input devices include push buttons, limit switches, temperature switches, and pressure switches. Input field devices have different operating voltage requirements, such as 120 VAC, 24 VDC, and 5 VDC. Discrete input modules with different voltage ratings are used to handle the different voltage requirements of these input devices.
- Regardless of the specific voltages involved, all input modules perform the same overall functions. They convert the input signals from the field devices to low-level DC signals the CPU can use. And, they electrically isolate the processor from the higher input voltages.

# Digital I/O

---

## ◆ Optical isolation (coupling)

- ❑ isolates the processor circuitry from the input circuitry
- ❑ the input module provides the CPU with a low-level DC signal which indicates the status of the input device.
- ❑ Any high voltage spikes which might occur at the input side will thus be prevented from reaching the processor and damaging it.



# Digital I/O

---

- ◆ The signal coming from the field input device is converted to a suitable level and filtered, then it passes through an optical isolation unit which consists of a light-emitting diode and a phototransistor.
- ◆ If the signal level is sufficient to represent an on input status (i.e. when the field device is providing a signal to the input interface), the LED will turn on.
- ◆ If the input signal voltage is too low, the LED will turn off.
- ◆ When the LED turns on, the phototransistor turns on. This on signal is provided to the processor, which records the status of the input bit as 1.
- ◆ When the LED turns off, the phototransistor will turn off. The status of the input bit will then be changed to 0.

# Digital I/O

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## ◆ Discrete Output Modules

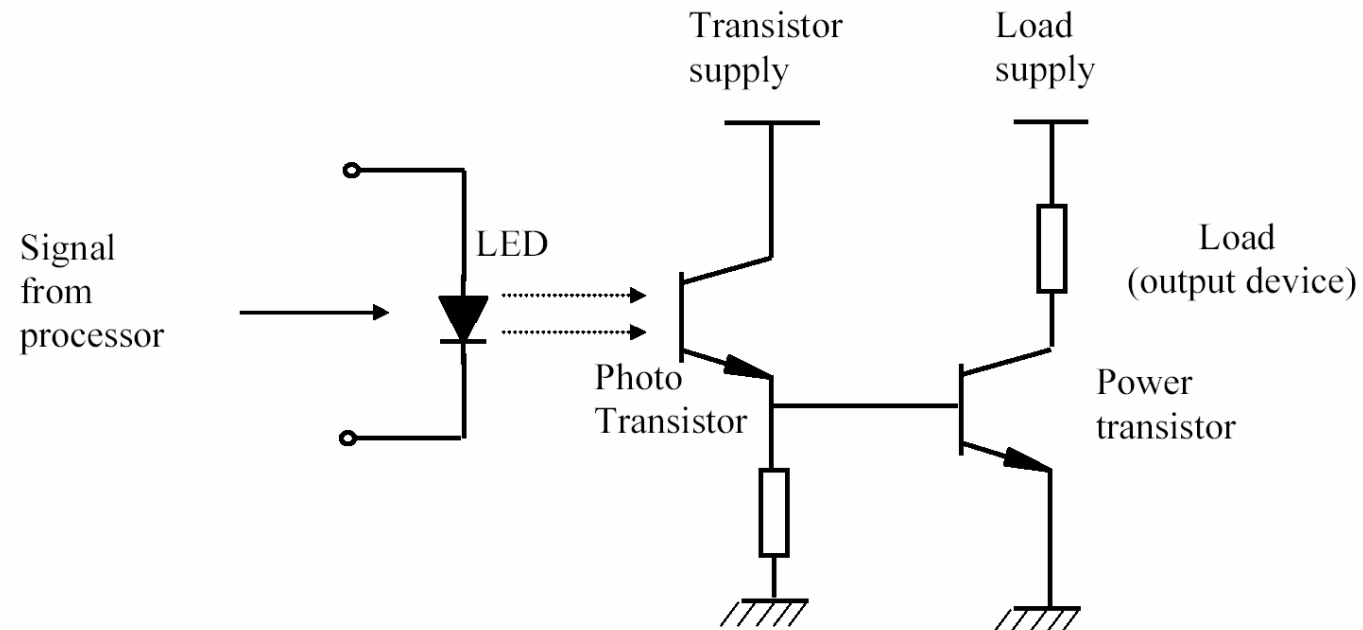
- ❑ One function of an output module is to convert the low-level DC output signals from the processor into appropriate voltage signals that switch the output devices on or off
- ❑ Typical discrete output devices include solenoids, motor starters, and alarm indicators
- ❑ Since different output devices have different voltage requirements, output modules are available with different AC and DC voltage ratings
- ❑ There are a variety of output modules. Two types that are commonly used are DC output modules and AC output modules. Only the DC output module will be discussed here as an example

# Digital I/O

---

## ◆ Discrete Output Modules

- ❑ A DC output module converts the output from the processor to the desired DC voltage signal that turns the output device on or off



# Digital I/O

---

- ◆ This circuit uses optical coupling to isolate the processor circuitry from the output device circuitry
- ◆ The power transistor acts as a switch, turning the output circuit on and off in response to signals from the processor
- ◆ For example, when the LED in the processor circuit receives an ON signal, it turns on
  - ❑ This causes the phototransistor to complete the circuit to the power transistor
  - ❑ This activates the power transistor and switches the output circuit on, signalling the output device to activate.
- ◆ When the signal in the processor circuit goes off, the LED goes off
  - ❑ This turns off the photo-transistor
  - ❑ As a result, the power transistor turns off, and the output circuit is de-energised.

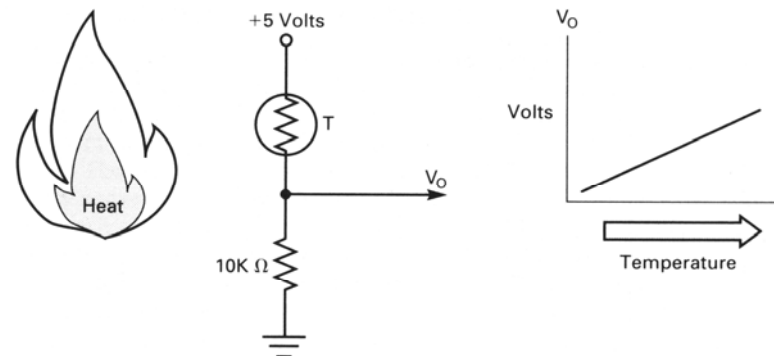
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# Analogue I/O

# Analogue Input

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- ◆ Most sensors produce an analogue electrical signal
  - ❑ DC voltage 0-10v, for example
  - ❑ DC current 4-20mA, for example
  - ❑ AC voltage, variable frequency
- ◆ Example: Thermistor temperature sensor
  - ❑ Resistance inversely proportional to temperature
  - ❑ 10k ohms at 25 degrees centigrade

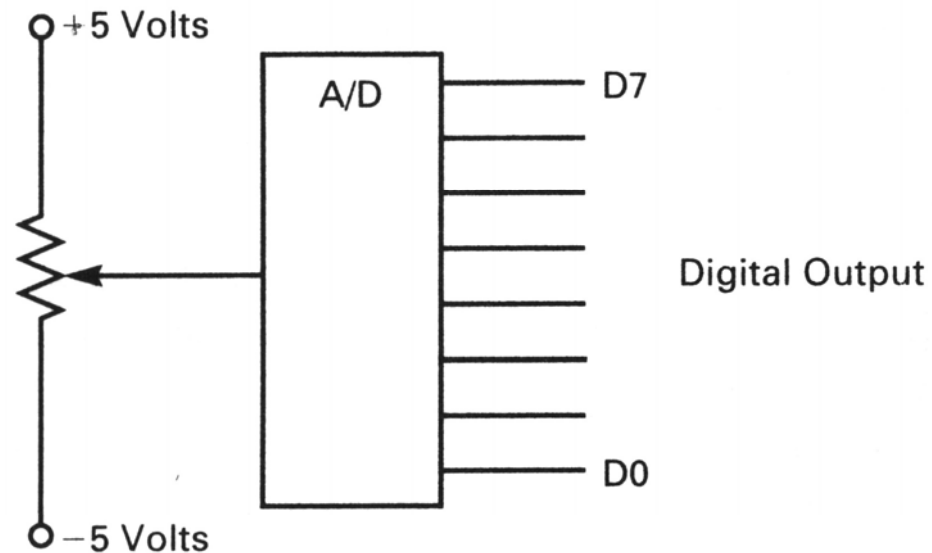


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# Analogue Output

8-Bit A/D Converter  
Input Voltage Range -5 to +5 Volts



<u>Voltage</u>	<u>Binary Output</u>	
-5	0000 0000	10 Volt Span
0	1000 0000	
+5	1111 1111	

# Analogue Input

---

## ◆ Analogue-to-Digital Conversion

- ❑ Converts analogue signals (mainly voltage) to a digital signal compatible with the TTL voltage required internally by the PC
- ❑ Resolution of ADCs depend on the number of bits used and the range of the analogue input. There are  $2^n - 1$  levels for an n-bit ADC:

$$\begin{aligned}\text{resolution} &= (V_{\max} - V_{\min}) / (2^n - 1) \\ &= (5 - (-5)) / (2^8 - 1) \\ &= 0.0392 \text{ volts}\end{aligned}$$

# Analogue Input

---

## ◆ Analogue-to-Digital Conversion

- ☐ It is important to choose an ADC with the proper resolution
- ☐ If the analogue voltage changes by an amount which is less than the resolution then it will not be detected, and the computer will not know about the change
- ☐ Cost is a major factor when selecting the resolution.

# Analogue Input

---

- ◆ Three types of Analogue-To-Digital Converters are commonly used in DA&C boards
  - ❑ **Dual-Slope converters:** simplest and least expensive (slow ... 10-100ms)
  - ❑ **Successive approximation converters:** most popular converter, faster than DS
  - ❑ **Flash converters:** very fast and expensive

# Analogue Input

---

## ◆ Sample and Hold Circuits

- ❑ Used to stabilise the changing analogue signal while the conversion process is taking place
- ❑ A typical technique is to include a capacitor which is charged to the value of the input analogue voltage

# Analogue Input

---

## ◆ Multiplexed-inputs

- ❑ This allows more than one channel (more than one analogue input) to be connected to the ADC
- ❑ Typically 8 to 16 channels are multiplexed on DA&C boards.

# Analogue Input

---

- ◆ In order to program any DA&C card, you need to know the register structure of the card and the base address
- ◆ Then you simply follow the manufacturer's programming procedure which involves writing to and reading from specific registers
- ◆ And you know how to do this in C (outportb and inportb).

# Analogue Output

---

## ◆ Digital-to-Analogue Conversion

- ❑ Convert a value represented by an discrete (integer) variable to a corresponding analogue quantity
- ❑ Again, resolution is important: there are  $2^n-1$  levels for an n-bit DAC; e.g. 8-bit DAC, 5V voltage range:

$$\begin{aligned}\text{resolution} &= (V_{\max} - V_{\min}) / (2^n - 1) \\ &= (5 - 0) / (2^8 - 1) \\ &= 0.0196 \text{ volts per step}\end{aligned}$$



# Analogue Output

---

## ◆ Digital-to-Analogue Conversion

- ☐ Programming a DAC is done in a similar way to ADC
- ☐ A data acquisition and control card typically contains both ADC and DAC
- ☐ Programming the DAC is achieved through writing to and reading from appropriate registers



LabJack USB AD&C: [www.labjack.com](http://www.labjack.com)

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